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THESIS

MICROLAN

FILE TRANSFER PROGRAM FOR MICROPROCESSORS

by

Roger Dean Jaskot
and
Harold Wayne Henry

March 1985

Thesis Advisor:

G.E. Latta

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MICROIAN
File Transfer Program
for Microprocessors

by

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requirements for the degree of

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from the

NAVAL POSTGRADUATE SCHOOL

March 1985

ABSTRACT

The age of automation has established its foothold in today's society. Computerization now affects almost everyone's job, and sharing of information is vital to successful job performance. Manual transfer of information is inefficient and prone to error, so another means is needed. One option is computer networking. Both Local Area Networks and long-haul networks presently exist, but they are either very expensive or hardware dependent.

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I. INTRODUCTION

The age of automation has established its foothold within the civilian as well as military communities. Very few jobs are left unaffected by computerization. Military or civilian, the "boss" is only as competitive as his or her information - which must be accessed or acquired from outside sources. Application of this to current automation implies sharing database information between computers. For example, consider two users of a manual, each holding part of that manual in their files. If one of the users suddenly needs information from the other user's portion of the manual there needs to be a means for access. In certain applications, especially military, speed is essential in information transfer. Transcribing data in the manual mode is ineffective due to slow response time and an increased chance of error.

One option that will help eliminate some of these problems is networking. There are two different types of networks presently in existence, long-haul and local area. Long-haul pertains to large geographic areas; examples being TELENET, TYMNET, ARPANET, and the public telephone system. Local area concerns itself with a much more restricted geographic area; examples being ETHERNET, OMNINET, PCNET, WANGNET, and LOCALNET 20. Our interest lies within the field of local area networks (LAN).

One major distinction between types of LANs is baseband versus broadband. Baseband is limited to transmission of bits of information while broadband allows transmission of video, audio, and digital data. Normally, baseband is also limited to single channel connections while broadband allows multiple channels for transmission. For our interests, the

key factor is that broadband requires expensive hardware and software. We have therefore focused on a baseband LAN system.

To focus our efforts even further, we compared asynchronous versus synchronous data transmission. In asynchronous transmission, machine interface is controlled by start and stop signals (handshaking) between the two microcomputers. Synchronous transmission requires both micros to operate on the exact same timing signals, either through a shared timing circuit or highly accurate timing systems at both ends. By themselves, no two computers - even of the same make and model - can be guaranteed to operate synchronously, and the cost of highly accurate timing is out of range for the small user. Since the military needs a low-cost, readily available system (see Chapter IV), we focused our efforts on asynchronous, baseband LANs.

LANs have become a common addition to many large organizations. They provide communication within a building or small groups of buildings, such as on a campus. Specific configurations depend on the volume and characteristics of the traffic, and the demands placed on the system. Local computer networks can also share peripherals. Sharing printers can be very cost-effective. By reducing the need of multiple printers, the idle time is kept to a minimum. Also, if one breaks down, the operator takes advantage of one of the other shared printers. Electronic mail may also justify the network depending on its implementation.

The military, which for our purposes is another large organization, has many of these same needs. The acquisition process varies between organizations. Standard acquisition methodology for the military is to evaluate present and future needs, come up with a list of requirements which a system could accomplish, competitively bid the system, and then await completion by a chosen manufacturer. This is an

over-simplified description of the actual process, but it will suffice for this discussion.

A large factor in determining which systems will be actually be acquired by the military is availability of funds. To alleviate some of the monetary problems, DoD tries to incorporate systems which can be used by the four major services; Navy, Air Force, Army, and Marine Corps. Lead-time required to obtain an operable network satisfactory to one or more services is usually measured in years, and the final product usually neglects the needs of some echelon levels.

To fill the time-to-acquisition gap, we have developed a program, MICROLAN¹ for transferring computer files between the types of microcomputers that are already present in the field. Since we use equipment and technology already available in the field, our miniature LAN can be quickly and cheaply installed. Cost is discussed in Chapter II.

To simplify the design of MICROLAN and ensure the flexibility of operating on different microcomputers, we did not include protection from collision of data on the transmission medium if more than one micro tries to transmit a file at the same time. If a second user tries to transmit a file while a file transfer is in progress, the receiver for the original file transfer will have a checksum mismatch with the original sender. Eventually the file will make it through, but transmission will be disrupted until the second transmitting unit decides to stop sending. As a result of this problem, a second file transfer session cannot be safely started until the first file transfer has been completed. We refer to this limited capacity of only one file transfer on the net at any given time as "low density" traffic.

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MICRCLAN allows transfer of computer files from a disk (hard or floppy) in one micro to a disk in another micro. The files can be man-readable data or text files or machine readable operation code. MICRCLAN also takes advantage of common equipment that is inherent to the majority of microprocessors. As an example, the RS232 is a standard interface connection on most microcomputers. Any medium that accepts the RS232 (hardwire, AC modem, phone modem, fiber optics, etc.) can be used with MICROLAN, whereas other file transfer programs are company/device dependent. A particularly cheap medium would be the use of existing power system wiring (ie. - the wall power plug) as an access to other computers. However, a device called an alternating current (AC) modem would be necessary to make such a connection. Such a device is now available off-the-shelf.

The training required to use MICROLAN is minimal. The necessary computer skills should already be present for those personnel presently working with the military systems. Actually, only routine clerical skills are needed for proper operation.

II. PROGRAM DESCRIPTION

The intention of this chapter is to give the reader a broad overview on the purpose and functions of MICROLAN. A detailed description involving assembly language is available in Appendix A.

MICROLAN is a file transfer program intended for use with most microcomputers. It's a very straightforward program designed to reduce the need for manual transcription and delivery of files. The reduction of error inherent with the manual transcription is a benefit with this system. The program can be used between microcomputers within an office or between different buildings. The design of the program is based on the lower networking levels (see Chapter 4). MICROLAN is readily available and can be implemented in any size military command or installation. It is presently written in assembly language for CP/M² - 80, CP/M - 86, and MS.DOS.³ (Copies of each program can be found in Appendices C, D, and E respectfully. Very minor changes will have to be made to the program, depending on which microprocessor is used). Using one of these three versions as a basis, MICROLAN could be translated to operate in another language, if needed. However, we feel that these three versions should be compatible with the majority of the systems presently operating in the military. MICROLAN has been tested on Northstar, Apple, and IBM microcomputers.

In MICROLAN, we have improved on asynchronous (character-by-character) transmission by adding the higher speed of synchronous transmission. Rosner states that:

²Registered trademark of Digital Research

³Registered Trademark of Microsoft Corporation

Low-speed, asynchronous character-by-character terminals operate in typical speed ranges of 75 to 600 bits/second. This class of terminal is a nonintelligent device and thus cannot respond to the protocol features of a packet switch interface. High-speed, synchronous block-by-block terminals operate in typical speed ranges of 1200 to 9600 bits/second. This class of terminal can range from non-intelligent - which can only respond to a very limited set of level 2, link-control commands - to highly intelligent, processor-controlled terminals - which can support all packet switched network protocol features with the possible exception of multiple simultaneous logical connections. [Ref. 1: p. 118]

MICROLAN is asynchronous in transmission method. However, due to its structure, MICROLAN can operate at the speed of synchronous transmission (theoretically, as high as 19,200 baud). The value of 19.2 kbaud is the practical limit for RS232 hardware units.

We realize that there are systems with faster transfer rates presently available, however, the hardware required increases the cost of the system dramatically.

Since MICROLAN is a baseband LAN, we will compare its cost to other baseband LANs. All of the costs given assume that the user already has the microcomputer to be used in the system. The cost for Ethernet is \$988 per user, with a minimum starting cost of \$2202 [Ref. 2: p. 151] Omninet costs \$650 per user with a \$2230 minimum [Ref. 2: p. 141] and PCNet costs \$742 per user with a \$1762 minimum [Ref. 2: p. 129] MICROLAN has no minimum cost for software (the program listings are in Appendices C, D, and E) or for a starter kit. The cost for hardware connections should be a maximum of \$50 for a small system, and for an AC modem connection would be about \$150 per user. By restricting MICROLAN's capabilities, we have been able to provide a readily available, low-cost LAN system. User interface with MICROLAN is minimal, and is driven by on screen instructions once the user has initiated program execution.

MICROLAN is intended for use by organizations as a convenience, and as an alternate means of sharing

information that is not time-sensitive. By "convenience", we mean that you can get or send the information without having to leave your work station. "Time-sensitive" means that the information loses value with every extra second that it takes to get to the receiver.

MICROLAN requires action on both sending and receiving ends to initiate transfer. The users must meet on the net at either a standard time (e.g., 0900 each Tuesday), or they must coordinate just prior to starting file transfer (e.g., a phone call saying meet me on the net and send the file). Timing as far as whether the sender or receiver starts first is not critical; however, the send portion of MICROLAN dies after about a minute with no contact. If this occurs, the computer must be rebooted. The receive portion of MICROLAN will wait indefinitely for contact from the sending micro.

To show how MICROLAN is used for file transfer, consider Capt X, who needs information from Lt Q on a new project. (Procedures would be the same if Lt Q needed a printcut of a file, but Capt X had the printer.) Assuming that the physical connections are already made, Capt X calls Lt Q and tells the Lt to send the file on the net in 15 minutes. At that time, both Capt X and Lt Q type "MICROLAN" followed by a carriage return. If they are using the CP/M - 86 or MS.DOS versions, the Capt and Lt will now be asked to select transfer baud rate from a menu (see BAUDMSG in Appendix F). The selected baud rates must match. Next, as the receiver, Capt X types an "R". The Capt is asked whether to write the file to the A, B, C (for CP/M-86 or MS.DOS), or default disk drive. Once the Capt selects the appropriate disk drive, his or her micro is in the receive mode and proceeds under control of MICROLAN until file transfer is completed. As the sender, Lt Q types an "S" to enter the send mode. The Lt is then directed to enter the name of the file to be transferred in the format "B:Filename.Filetype" where B

represents the B disk drive. If the file was in the C drive, the Lt would replace the B with a C. If Lt Q typed in the filename in the format "Filename.Filetype", MICROLAN would assume that the file is on the default disk drive. Once Lt Q has entered the filename, MICROLAN takes over and no further action is required of either user unless Lt Q decides for some reason to abort file transfer. If the Lt should decide to do so, he or she could stop sending the file by pressing the <Control> and "C" keys at the same time just after a "*" has been printed on the screen to indicate that a 128-byte frame has been acknowledged.

MICROLAN begins by sending and receiving "handshaking" indicators allowing the micros to become synchronized. After the program is satisfied that they're in sync, the transmitting micro sends the filename and ensures through error checking that the correct filename was received. After this acknowledgement, the transmitting micro begins sending 128-byte blocks of information across the line. A checksum is calculated throughout transmission and is checked after each block is sent. If it checks good, then transmission is continued with the next block. If an error is detected, then the block is retransmitted. This procedure is continued until the entire file is sent. When the end of the file is reached, an "end-of-file" indicator is sent, telling the receiving micro that no further file information will be coming and to go ahead and close the file. A handshaking process then takes place, acknowledging file transfer is complete, and that both micros are ready to return to the operating system. Both micros then exit the program and are ready for the operators next desired command. It must be noted that, while MICROLAN is executing, the two microcomputers cannot be used to perform any other operations.

There are some safety factors incorporated for ease of operation. First, if the transmitting operator decides to abort transmission at any time, an input of "control C" will execute the abort. A message will let the receiving operator know that file transmission was aborted and that an empty file exists under that filename. Second, if the receiving file already has an existing file with the same filename, or if the transmitting micro cannot find the desired file, execution will stop, advising both sides of the situation. Third, if the receiving micro has a full disk and cannot receive the entire file, transfer will be aborted and both the operators will be advised. Appendix A expands on the above routines if any clarification is needed.

MICROCLAN has been tested for operation in sending both man-readable text/data files and machine-language command files. It has been tested for transfer at 1200 baud between two Apple micros, two Northstar micros, Apple to Northstar and Northstar to Apple. Tests were also run at 4800 baud from Northstar to IBM PC, IBM PC to Northstar, and between two IBM PCs. Operation of MICROPLAN was also tested at 9600 baud between two IBM PCs; however, the code logic used in MICROPLAN proved unable to cope with the timing problems at this speed. Future revisions could overcome these problems. These tests were performed using hardwire connections; however, results should not vary with different connection media.

III. MILITARY APPLICATIONS

MICROLAN can be used by the military to help meet current information sharing needs. Its attributes help alleviate certain problems that exist in current systems or that arise when obtaining a new system. One major problem the military encounters is the time delay that exists between statement of need and delivery to the service. Too often is the case that when the finished product is finally fully operational, the "threat" is at an advanced stage, thus making the system somewhat obsolete. Rather than waiting for a technological breakthrough to occur that will take care of any possible future threat, a system has to be deployed to counter the current "threat". MICROLAN is available for immediate implementation. It can be used by itself, as an enhancement to existing systems, or in the development of future systems.

The option exists for intra- as well as inter-service use. If the need for joint interoperability doesn't arise in a specific situation, MICROLAN is still fully operational within the realm of a single service. Inter-service use poses no major changes either. The same existing hardware and software are still used. MICROLAN can be quickly adapted to almost any microcomputer, thus overcoming the profusion of dissimilar equipment in the field.

Required operational training of personnel is kept to a minimum with MICROLAN. The military employs a vast range of users, varying in educational backgrounds. MICROLAN's ease of operation is limited only by the most rudimentary knowledge of the typewriter keyboard.

Cost overruns, scheduling delays, contract disputes, and a myriad of other pitfalls plague the Department of Defense

budget. MICROLAN is inexpensive, available, and easy to incorporate. Adhering to these attributes, MICROLAN would not be a financial burden to the military. "Word of mouth" is one of the best promoters of a new product. Enhanced by promotional meetings, computer bulletin boards, satisfied users, etc., MICROLAN's usefulness will hopefully be widely disseminated to all facets of the military.

This chapter presents some examples of how MICROLAN could be used by the military. It will also cite some examples of how the military is presently trying to automate information distribution.

A. NAVY

The U.S. Navy has undergone a major face-lift over the past two decades. Significant breakthroughs in technology has offered tremendous advances in ship-building design and associated weapon systems. Due to these advancements, decision-making by the warfare commander has been quickened by shorter planning cycles, dissemination of orders, and resulting outcomes of those orders. The "real-time" response to any attack has had to be critically shortened in order for present day operations to be successful. Turn-around time in paperwork has also gone through many changes in attempts to minimize slack time caused by tedious, but necessary, record keeping. Personal files, parts orders, and safety statistics are just a few of the necessary information requirements for any large organization.

Whether it be in an operational setting, such as the Combat Information Center (CIC) aboard a ship, or in a shore-based supply facility, the Navy is always looking for ways of reducing the workload placed on its personnel. One such system that the Navy is presently pursuing is the ZOG

[Ref. 3] system, which has been placed on board the aircraft carrier, the USS Carl Vinson, as one of three microcomputer networks. This example will provide an idea as to what the Navy is looking for in the field of computers.

ZOG is a general-purpose human-computer interface system that combines the features of a database system, a word processing system, and an operating system shell. This system is a distributed database system implemented on a network of 28 high-powered personal computers (PERQS), interconnected via a wideband local area network (Ethernet).

The uses of a local area network with computers are seemingly endless. A few examples of the ZOG system will suffice. On board the USS Carl Vinson, ZOG has been used as a software management database, well suited for structured software development. It has also been extensively used to implement forms of electronic communication, such as electronic mail, bulletin boards, and teleconferencing. In a more advanced area, ZOG was used for project management; to develop multi-level task structures which could be used not only for planning, but for implementing and evaluating as well. Other areas that were explored were training, interfacing with an existing system, and retrieval of emergency operating instructions (in this case, for commercial nuclear power plants). As with almost any new system, there's always room for improvement. An extension of ZOG is the Knowledge Management System (KMS). In KMS the model of a frame has been extended to include graphical as well as textual items.

The ZOG example provided a good insight as to where the Navy is looking in terms of newer technologies. Akscyn and McCracken brought out a good point in their report (Ref. 2). That is, how the users of the system can make their work usable by others, especially since there are few situations in the real world where people do not depend on interaction with others to accomplish their work.

Our file transfer program, integrated in a local area network, could alleviate some of the problems. To gain a better perspective on the usefulness of this program, let us state that this project was not intended to be used in a time-sensitive environment. An example of that would be in use with the Navy Tactical Data System (NTDS) updating friendly as well as enemy positions. In this situation, seconds are critical concerning command decisions.

One area where this system could be very useful is in the supply system. It is irrelevant as to whether the supply department involved is shore-based or afloat. Transferring files, part orders, etc., between buildings or ship compartments would drastically reduce the manual labor presently involved. Consolidating the payroll system would greatly reduce the space required for all of the necessary paperwork.

Electronic mail would be a good use also. The administration departments would find it useful in preparing command-wide bulletins (e.g. Plan-of the Day) or collating fitness reports. The communications department could utilize the system for drafting message traffic. Instead of congesting the commanding officer's desk with messages awaiting approval, they could be sent to his disk, which he could then address at his own leisure, returning finished copies at will. The maintenance department could "converse" with the supply department in a more organized manner concerning needed equipment. The safety department, in conjunction with the maintenance department, would be able to pass or collect necessary statistics needed for periodic reports.

These are just a few examples which could be incorporated within a command. They would not have to utilize these opportunities all the time, however the option would be there. The main benefit of this system is elimination of

transferral of paperwork between departments (or even within departments). Having a condensed file of needed information on one disk would definitely reduce the amount of lost information due to scattered, and inadvertently discarded, paperwork. One important aspect to keep in mind is that the manual method of information transferral would still be available, if needed for one reason or another.

B. ARMY

The U.S. Army does not enjoy the luxury of being numerically superior to present day opposing forces. Even though the Army has a slight qualitative and technological advantage, the threat combines its numerical advantage with its increasing weapon and combat technologies to at least nullify the slim margin the U.S. presently holds.

The Army, like the other services, tries to utilize as much new technology as possible to sustain this margin. There is more information on and about the battlefield today than ever before, however, the staff essentially still processes the information in a manual mode. There are some automated procedures, but the bulk of the system contains mostly manual procedures.

In order to alleviate some of these problems, the Army has introduced CPASS (Command Post Automated Staff Support System). [Ref. 4] The primary purpose of the CPASS system is to provide automated assistance in performing staff functions. The automation devices and software of the system are tools that expand the staff's capability to handle more information and to utilize the information more efficiently. Some of the intended uses for CPASS are:

- a) An information processing system to develop and execute staff plans and operational orders.

- b) Provide a near-term staff wide command post automated information distribution and decision support capability.
- c) Provide more real time and near real-time accurate information to commanders and their staffs.
- d) A graphical situation display and hard copy overlay capability.
- e) Automation for both tactical and garrison applications without a requirement for intensive train-up or transition to meet deployment or operational requirements. Such a system is required for daily use, not just to support the deployed command post.
- f) A initial capability for the evolutionary development of concepts, doctrine, procedures, hardware, and software for the continuing automation of command post staff activities.
- g) Capability to support the dispersed command post in accordance with current doctrine. It must demonstrate the additional operational and organization changes required to support the dispersed command post.

Items b,e,f, and g are along the same intentions (automation of commands, information sharing, and minimal training requirement) as that of MICROLAN. The hardware/software make-up of the CPASS is unquestionably larger and more complex than our system. However, some of their components and structures are used for the same basic purposes as ours. A few of them are:

- a) Within the command post cluster, devices are interconnected by a local area network (LAN). The LAN is physically versatile and can interconnect devices in all shelter configurations of a command post cluster (expandable shelters, buildings, or within an armored command post vehicle or van). Media types to be used in LAN include twisted pair wire, coax cable, or fiber optics.

- b) A Network Computer Unit (NCU) containing a network interface element and a microprocessor performs data routing functions within the work station, and between the work station and other cluster devices.
- c) The file storage device stores elements of the data bases of other command post clusters in anticipation of combat loss or equipment malfunction.
- d) The communications processor, in conjunction with the communications devices of the area communications system, provide end-to-end message transport service to allow essential interstation communication, such as message routing, data base interactions and graphics data transfer.

The required personnel training for CPASS is very similar to ours. The introduction of CPASS is not projected to require increased command post manning levels, new military occupational specialties, or Army skill indicators. Operators are those who are already assigned to command post staff functions. Routine clerical skills are the minimum essential personnel qualification skills needed to operate the system (i.e., typewriter keyboard, filing, etc.).

The CPASS example is a good indicator of what direction the Army is heading in terms of computer use. Our system contains many of the same qualities that the Army desires and requires.

C. AIR FORCE

The Air Force is deeply involved in networking headquarters and tactical functions. However, the focus is on expensive broadband networks and tends to neglect the smaller users. The report on the Hardened Tactical Air Control Center (HTACC) even states that:

Of the three major functions in the HTACC--intelligence, operations, and logistics--the direct automation support from the CONSTANT WATCH Program is largely restricted to the intelligence and operations functions under current plans. Indirect support for logistics functions will result from the automation of intelligence and operations functions which logistics uses and from secondary use of the communication capabilities implemented under the CONSTANT WATCH program. Eventually, logistics automation requirements must be addressed, and, hopefully, integrated with the intelligence and operations activities. [Ref. 5: p. II.4]

The HTACC involves use of high-cost broadband technology. Our method could use the power cables that must be present anyway, and requires only an AC modem and minimal cabling in addition to the RS232 that is standard on most microcomputers. Since the logistics requirements are not real-time, and shouldn't be extremely high density in traffic, they could be supported by MICROLAN. Reports on status of supplies and requests for movement of supplies are the types of traffic that could be expected on the system. Also, minor information transfer between control positions in the TACC could be accomplished on our system, reducing the workload on the real-time LAN system.

The Air force is also working on a LAN (PENTANET) for the Pentagon to provide:

- a) The exchange of data between local and remote interactive Keyboard Video Display (KVD) terminals and local and remote processors, wherein local and remote connote devices within and exterior to the Pentagon
- b) The electronic exchange of variably formatted reports and documents between local and remote workstations
- c) The local and remote distribution of digitally encoded graphic and facsimile products
- d) The transfer of data files between local and remote processors and between local and remote peripheral devices

- e) The transfer and distribution of teleconferencing and commercial video and associated analog voice, and low speed analog and digital control signals
- f) The switching and exchange of analog and digitized voice signals between users. [Ref. 6: p. 9]

Here, our program could be used to supplement the PENTANET as an interoffice message transfer system, reducing the major net's workload. The remos could include coordination on letters or short documents that can also be sent using our program. File transfer via MICROLAN is not limited to text. If one office has a program that another would like to use, it can be passed over MICROLAN, even if the computers are not the same brand name product.

Another Air Force function that could make use of our program is the Base Information Transfer System (BITS). BITS is the base mail system. Memos, completed forms, blank forms, appointment notifications, general mail, and coordination copies of documents are transported around base using this system. Base administrative personnel pick up the correspondence, take it to a central processing office, and then deliver it to the destination. From past experience, BITS has been known to lose messages, and the only way that has been recommended for improving timeliness of service is an increased number of delivery runs [Ref. 7]. To implement more runs would require more personnel, therefore increasing costs. Use of MICROLAN would have a low one-time cost and almost no upkeep. One additional requirement for implementing our system on a base-wide basis, if power system wiring is to be used as the net, would be installation of capacitors on power transformers to allow the LAN to cover a wider area of the base. The capacitors would allow our signal to pass through the transformer while preventing the AC power from crossing over. Installation would be only a minor problem. Memos and coordination would require no

modification to be passed using MICROLAN. Reports and supply requests could be formatted and transferred to action agencies for printout on the receiving end.

One possible use was suggested by Hq. Tactical Air Command, Tactical Air Forces Interoperability Group (TAFIG). TAFIG identified a need for transferring wing or squadron databases to computers onboard aircraft. This would require either a disk system in the aircraft or storage of the program on a programmable memory chip, which would be more reasonable. This system would provide pilots with data through the onboard computer, decreasing some of the time that would be required for briefings on the ground.

D. MARINE CORPS

We contacted the Marine Corps Command and Control Systems Office at Camp Pendleton, California to determine what types of network systems they were looking for. They indicated that they have immediate needs for interoffice file transfer and mailgram systems, both of which MICROLAN can provide. They also have a need for a tactical message transfer system in the 9600 baud transfer range, which would have to be able to be sent via encryption or other secure means. Our system of transfer using RS232 technology and the buffers built into MICROLAN should allow transmission via a variety of media, including fiber optics. Although we have not tested MICROLAN with encryption, we do not expect serious problems in doing so. In addition to these uses, the CPASS system mentioned under the Army section of this chapter contains uses that would also apply to Marine operations.

E. CHAPTER SUMMARY

In this chapter, we have presented several possible uses for our file transfer program in filling present requirements of the four service branches. We have not attempted to enumerate every possible application of our program, only some representative uses for each service. There are most certainly more file transfer uses that exist that MICROLAN can be applied to. The key prerequisites for using our system are that the data is not time sensitive and that traffic is low-density. A review of our suggested uses for MICROLAN shows that there are applications throughout the spectrum of service organizations-whether in the back office or on the battlefield, shipboard or aboard aircraft-that meet these prerequisites.

Use of MICROLAN for interoffice memos could be applied to any installation or organization. For example, the Navy Postgraduate School has a need for an inter- and intra-departmental mailgram system. Intra-departmental networking should be no problem for an AC modem system, since members of a given department are usually grouped together in the same building. For inter-departmental use or departments that are spread across campus, capacitors would have to be used as mentioned in the Air Force section of this chapter. If use of the system becomes saturated, methods identified in the Conclusion for separating nets could be employed.

It is important to note that the use of RS232 interface technology allows a varied means of connection between sending and receiving units. This is a significant factor in MICROLAN's flexibility. Another aspect of MICROLAN that contributes to its flexibility and interoperability is that it is confined to the lower levels of the International Standards Organization (ISO) model. This ensures that neither higher levels of computing power nor

specialized components or exotic software are required to implement the MICROLAN system. Chapter 4 explains how MICROLAN fits into the ISO model.

IV. MICROLAN AND THE NETWORK MODEL

To understand where MICROLAN fits into the International Standards Organization (ISO) Open Systems Interconnection (OSI) model, the basic basic idea of that model's concepts are given. The ISC OSI model consists of seven layers (levels) corresponding to computer functions and interconnection. These range from a basic physical layer to user interface. Figure 4.1 is the standard representation of these layers.

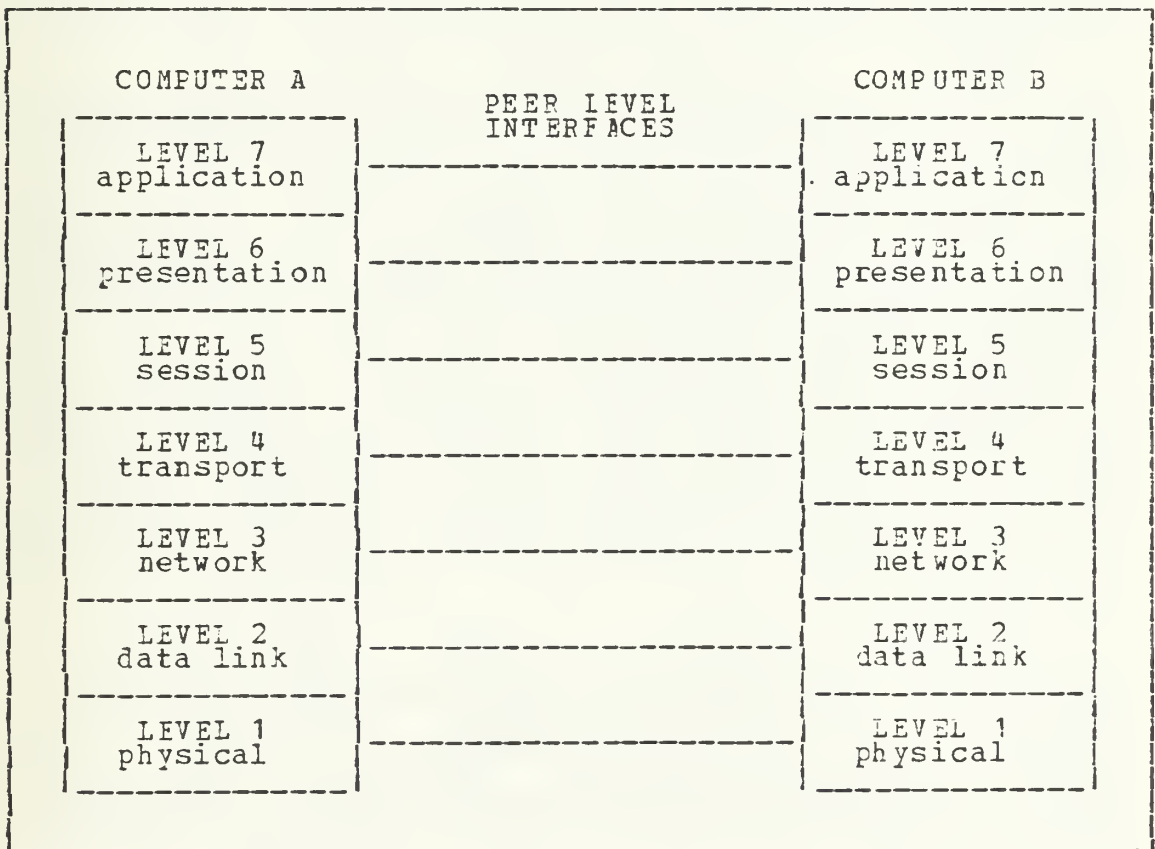


Figure 4.1 International Standards Organization Protocol Model.

A. MODEL OVERVIEW

The seven layers of the OSI model are discussed in short and at length by many authors on networking computers. The best summary that we found was Roy Rosner's book in which he states:

The lowest level of the ISO protocol hierarchy is the physical level, where previously defined standards were applied to define the physical interface. By physical interface to the network we refer to the pin connections, electrical voltage levels, and signal formats. Level 2, known as the data-link level, controls the data link between the user and the network. This level defines data format, error control and recovery procedures, data transparency, and implementation of certain command sequences. For nonswitched networks, or the interface of simple terminals with computers through point-to-point services, generally only levels 1 and 2 are required. Networks designed by a single manufacturer around a single product line, generally do so with a combination of level 1 and level 2 protocols.

Level 3, the network level, defines most of the protocol-driven functions of the packet network interface, or the internal network. It is at this level that the flow-control procedures are employed and where switched services are initiated through a data call establishment procedure.

Level 4, known as the transport level, assures the end-to-end flow of complete messages. If the network requires that messages be broken down into segments or packets at the interface, the transport level assures that the message segmentation takes place and that the message is properly delivered.

Level 5, the session control level, controls the interaction of user software, which is exchanging data at each end of the network. Session control includes such things as network log-on, user authentication, and the allocation of ADP resources within user equipment. Level 6, the presentation level, controls display formats, data code conversion, and information going to and from peripheral storage devices. Level 7, the user process or user application level, deals directly with the software application programs that interact through the network.

Although at levels 5, 6, and 7 the protocol is defined from a functional viewpoint, implementation of standard software that can operate at these levels has been slow. The software at all of these levels (often referred to as peer-level software) tends to be both equipment and application dependent. However, the layered approach to protocol development achieves a degree of isolation and modularity between the various layers, so that changes in one level can be made without changes in any other level. [Ref. 1: p. 109]

MICROLAN's structure fits into the lower levels of the OSI model. For our purposes, it is important to note that layer 1 signaling modes include: full duplex, half-duplex, synchronous, asynchronous, balanced, etc. There are also several standards that exist at layer 1. For example, there is EIA's RS232 and RS449, and CCITT's X.21, V.24, and V.35. [Ref. 6: p. 97] How MICROLAN functions in each layer is explained in the following pages.

B. PHYSICAL LAYER

For the Layer 1 interface, we take advantage of RS232 technology, thus providing a standardized physical interface for MICROLAN. This eliminates the problem of matching high and low voltages for different computers. Normally each individual bit is regarded as an entity for Physical Layer purposes; however, in our design, an 8-bit byte is used as an entity for transmission of data. This is the smallest segment of information handled by a microcomputer's accumulator, and is the ASCII representation of data characters. It is in this layer where we get our greatest flexibility. This flexibility arises from the fact that a variety of methods exists for linking one RS232 to another, as mentioned in Chapter I, providing the user with options in the type of hardware they can use.

C. DATA LINK LAYER

In this layer, our data is grouped into a 'frame' of 128 bytes. This number equals the storage capacity of the Direct Memory Access (DMA) buffer that is standard on microcomputers. A file is broken into frames and reassembled using the microcomputer's operating system commands. On the sending side, a read sequential command breaks out the sequential frames by reading 128-byte blocks into the DMA

for transmission. On the receiving end, the DMA is filled by the 128 bytes that were sent, then a write sequential command places the frame into the new file in sequence. By doing this, we prevent having to develop software methods for sequencing frames.

MICROLAN performs error checking on a frame to frame basis. Within each frame of data, a checksum is calculated by both sender and receiver and compared at the receiving end. If the two checksums don't match, the receiving micro informs the sending micro, which then retransmits the same frame of data, repeating until the frame is acknowledged as received correct or the sending user decides to abort file transfer. Combined with the error checking, we built buffers in to allow slower micros (e.g., Apple versus Northstar) to catch up to their faster counterparts.

Since we use the DMA regulated 128-byte block for our data frame on both ends of transmission, the amount of data sent at one time can never exceed the receiving micro's buffer capacity. Therefore, MICROLAN doesn't require a 'buffer space left' notification that would normally occur in this layer. [Ref. 8: p. 17]. Instead, it is at this level where the receiving micro checks for free disk space and informs the sending micro to abort file transfer if there is no more disk space for storage. Finally, one frame must be acknowledged as received and correct before MICROLAN will send the next frame. This eliminates the problem of duplicate or lost data frames.

As stated by Rosner in the above quote, this is the highest level required of simple, nonswitched networks. However, in order to allow the user some control of MICROLAN, we do provide some features in the Session layer.

D. SESSION LAYER

This is the final layer used in MICROLAN. Here, the user invokes MICROLAN and initiates connection by selecting send or receive functions. During this process, the user also selects which disk drive (default or an alternate) for accessing or storing the file. The sending user's option of aborting file transfer also falls under the definition of this layer. See Chapter II for operation instructions.

E. SUMMARY

The majority of MICROLAN's activities occur in the lower two layers of the ISO OSI model, as seen above. As a result, user friendliness is limited. Also, as mentioned in Chapter II, MICROLAN monopolizes the computer, allowing no other operations. MICROLAN is being used as the basis for a higher level network system by a fellow NPS student, LCDR Jeanie Egbert, in her thesis, "A MICROCOMPUTER NETWORK: INVESTIGATION AND IMPLEMENTATION". The combination of MICROLAN and her thesis provides a more presentation oriented structure. LCDR Egbert's LAN system allows the user to perform other operations on their micros while files are being transferred.

Since MICROLAN performs no Network Layer functions, no collision detection or prevention is provided - as mentioned in Chapter I. Flow control is limited to the link between one sending and one receiving micro on the network.

By limiting MICROLAN's main functions to the lower layers of the ISO OSI model, we have provided a simple, nonswitched file transfer system (LAN). MICROLAN is designed to operate on a variety of microcomputers, not just one product line. Therefore, we have gone one step farther than Rosner indicated in the first paragraph of his model description.

Further improvements can always be made to a program.
We identify some of these possibilities in the next chapter.

V. CONCLUSIONS AND RECOMMENDATIONS

A. CONCLUSIONS

The need exists for a public domain low-cost file transfer system to provide an alternative to commercial systems (e.g., ETHERNET). The type of transfer considered here is low-density traffic that is not time-sensitive. Types of files include computer programs (both in text and in machine code), data, messages, and text files.

Our solution was to take advantage of standard RS232 technology that is used by all microcomputers. This makes MICROLAN capable of operating on a wide variety of micros. Writing MICROLAN in CP/M, CP/M-86, and MS-DOS versions of assembly language makes it compatible with-at a minimum-Northstar, Apple, and IBM-PC compatibles.

There are two key restrictions encountered when using MICROLAN for file transfer. First, MICROLAN monopolizes both sending and receiving micros so that they are not available for other purposes until file transfer is completed. LCDR Egbert's thesis, mentioned in Chapter IV, addresses this problem. Second, only one file transfer can be conducted on a particular net at a given time. This is because MICROLAN has no detection or prevention of on-line data collisions. Providing multiple paths in a given network area would reduce chance of collision and allow more than one transfer at a given time. A method for doing this is discussed in the second part of this chapter.

Our intentions were to keep MICROLAN simple, so that an in depth knowledge of computers is not required to use it. Once the RS232 connections are made (standard plug connectors make this step relatively simple), entering 'MICROLAN'

from the console presents the user with easy to follow instructions including sample inputs that lead to file transfer. During execution, MICROLAN prints phrases on the user's screen, informing them of transfer status. These on-screen comments are shown in Appendix F. Also, a '*' is printed on the screens of both sender and receiver for each 128-byte block sent. A 'b' is printed on the sender's screen for each unmatched checksum. This feedback allows the users to see that transfer is occurring and determine if problems exist (e.g., several b's in a row). The sending user can exercise the option to abort transfer if the bad checksums persist.

Two safety features are included in the SLAVE subprogram. If the receiving disk already has a file with the same name as that of the file being transferred, file transfer is aborted and the users are informed of the problem. This protects an existing file from being overwritten by a new file. The other safety feature is activated when the receiving disk runs out of memory space. When this occurs, file transfer is aborted and both users are informed. The file is not closed under this method, so file transfer is 'all or none'. We chose this method because, especially in the case of command files, serious problems could arise from partial transfer of files. Command files cannot be readily repaired by the receiver by merely 'typing in' the missing parts.

In Chapter III, we discussed various military uses for MICROLAN, based on a search of network requirements from the four services. The main criteria for using MICROLAN is that the information requirement must not be time sensitive and traffic must be low-density. MICROLAN would also allow organizations that can't justify the expense of a high-powered network, an option for a low-power, medium speed (up to 19,200 baud) network at a much lower cost. Use of RS232

standards means that the user can take advantage of whatever connection medium is readily available, whether it be telephone, power lines, or direct wire. This would also help lower costs.

MICRCLAN operates mainly in Layers 1 and 2 of the ISO OSI network model, as explained in Chapter IV. It is a simple, transfer-oriented system. User interface at the upper levels is the minimum necessary to operate the program. This was done deliberately to maintain maximum flexibility in rewriting MICROLAN to run on different micros.

MICRCLAN has been successfully tested for operation using hardwire connections between the RS232s, at rates up to 4800 baud. When used with interrupt driven programs such as LCDR Egbert's, where timing problems will not exist, we expect transfer speeds of 19,200 baud to be possible. Replacing the hardwire connection with modems, fiber optics, or any other type of medium should not affect operation of MICRCLAN.

In MICROLAN, we have provided a flexible, low-cost mini LAN as an new option for information transfer. Of course, improvements can always be made to any program, so the next section presents some that we recommend for MICROLAN.

B. RECOMMENDATIONS

As mentioned earlier, MICRCLAN does not provide collision detection and prevention. One project for further research would be to develop program code to incorporate collision detection and avoidance into our program.

One change that would only require minor modifications is to return the user to the Send/Receive/Exit menu after file transfer is completed. We felt that returning the user to the operating system was more appropriate, but others may

feel differently. At this same level-i.e., the menu-it would also be possible to add the ability to send more than one file in a session. This would require changes in the File Control Block load subroutine, as well as a change in the end-of-file subroutine to loop back to send the next file.

It is possible that noisy transmission lines could cause problems with MICROLAN's checksum procedure. A subject for further research would be development of an algorithm for noisy line error checking perhaps by using cyclic redundancy.

As presently written, MICROLAN dumps files only to a disk system. To add flexibility, menu driven subroutines could be added to allow file transfer directly to other peripherals. This would allow one user to 'borrow' another's printer without moving it. Of course, file transfer would be slowed by the limited speed of the printer.

Use of MICROLAN as a Bulletin Board system would require a menu item in addition to Send/Receive/Exit. Subroutines to execute this option would have to use the Console Buffer and Random Access Memory of the micro to store bulletin items. The first bulletins would print on the user's screen, with following items stored in memory. The system would have to allow the user to page through the bulletin items using console keys. The option to send as well as receive items while retaining the previous items would also be helpful.

To allow up to 500 micros to communicate in a given area, the net can be broken into separate subnets. Each subnet would operate on a different frequency channel as set up by a central controller. In the example of the NPS net requirement, one channel could be for the Superintendent, one for logistics, one for each of the departments, etc. In

a one megahertz band, there could be 10-20 channels depending on baud rate. Since this is hardware driven, no software change would be necessary. However, channel selection could conceivably be software driven. If users are on several nets, they could use scanners to 'listen' for messages on the different nets in the same manner as radio scanners are used to listen for messages on Citizen's Band frequencies.

Tied in with 'listening' for messages, subroutines could be added to allow each user to have a personal identification number (PIN), assigned by net control. The micro would listen for messages to all users or to their PIN specifically, ignoring all others. This would operate best in conjunction with higher level programming, such as LCDR Egbert's program, that would allow the user to perform other computer operations while MICROLAN is looking for messages.

Our final recommendation is one that would make MICROLAN operate as a token ring network. On board ship, where power is not a problem, the micros could be left on continuously (actually this is better for the micro). In conjunction with the PIN idea, software changes would have to be developed that would allow MICROLAN to be used as an intercom system. One user would control the intercom, passing control to other users as they have the need to ask or answer questions. Control of the intercom would then be passed back to the master user.

Obviously, we have not covered every possible use or improvement for MICROLAN, but we hope that our description of MICROLAN and its possible uses has planted a seed for future research and expansion of low cost LAMs.

APPENDIX A

DETAILED PROGRAM DESCRIPTION

Prior to writing assembly language code for MICROLAN, we developed a flow diagram to show what we wanted to accomplish with the program. We developed the Master and Slave portions in parallel, showing rendezvous points with connecting lines. This flow diagram is shown in Appendix B. From this flow diagram, we developed subroutines to actually execute the steps and loops required to transfer a file. The programs shown in Appendices C, D, and E include additions that make MICROLAN more user friendly (e.g., ability to select which disk drive or to abort transfer).

The MICROLAN file transfer program consists of two subprograms that operate on separate micros, with frequent rendezvous to ensure parallel operation. We used modular programming style and developed the Master and Slave subprograms in parallel to insure that the two would rendezvous at matching subroutines. Data transfer is up to 8 bits per byte (ASCII or standard hex). Buffers had to be added in the Master rendezvous subroutines to allow the Slave subprogram to catch up when using different micros.

Our program is written to be used on microcomputers using either CPM, CPM86, or MS.DOS computer program/manager operating system. To allow use on other types of systems, changes will be required in the assembly language code to match that used by the micro to be used.

The remainder of this description refers to the CP/M-80 version of MICROLAN except as noted. First, MICROLAN must know which language format will be used during operation. Language format refers to the type of commands inherent to the microcomputers system. For example, the Apple loads

data from the input buffer into a memory address before reading it to the accumulator, while the Northstar takes data directly from an input port to the accumulator.

If the micro operates like a Northstar, the main change needed is in the definitions at the end of the program code. The user must change DATA EQU 04H to the number of the micro's input port and STATUS EQU 05H to the number of the status port. If the micro operates like an Apple, DATA1 and STATUS1 must be changed to reflect the micro's correct port numbers. The user should also verify that TXRDY and RXRDY reflect the correct values for their micro. Since the program is matched to Apple, Northstar, and IBM (and compatibles) types of microcomputers, MICROLAN should be useable by a wide variety of systems. There are seven subroutines affected by changing micros. They are POUT, STATIN1, STATIN2 and PIN. Slave also has the subroutine PIN1 that is affected, and Master has STOPS and GOCPM that are affected. They're easily spotted because they are the only subroutines that use the IF statement. Also, the appropriate constants will have to be added at the end of the program.

At the beginning of the program (see Appx C), the name of the micro that you are using must be set equal to TRUE. For example, APPLE EQU TRUE. The name(s) of other micros must be set to NOT TRUE. For example, NORTHSTAR EQU NOT TRUE. This activates the appropriate portion of the IF-THEN statements.

You will notice that we set the origin of MICROLAN at 0100 Hexadecimal (Hex). This is the standard position to load a program for execution. We then move the Stack Pointer to a higher memory address to prevent it from being overwritten.

To invoke MICROLAN, type 'MICROLAN' and press return. At this point in the program, if the user is using CP/M-86

or MS.DOS, they are asked to select transfer baud rate from a menu. For the CP/M-80 version, and continuing for the other two, the next step is the INIT subroutine asks you if you wish to send or receive. The HOLDING subroutine locks for a keyboard input until either an 'S', an 'R', or an 'X' is found. An 'S' sends the program to the MASTER subprogram. An 'R' sends it to the SLAVE subprogram. An 'X' returns the user to the main operating system of the micro.

MASTER first asks for the name of the file to be sent. The user can also identify at this point which disk drive to retrieve the file from. Possible selections are A, B, (for CP/M -86 and MS.DOS systems, the option for a C disk is also included) or default. The user specifies the drive by typing in the format 'B:filename.filetype'. If no drive is specified, the default is assumed. While the user is entering the filename, the FILUP subroutines prepare the File Control Block (FCB) for receiving drive information and the filename. The FCB is located starting at memory location 005C Hex and is 32 memory locations long. It is the default filename location for all microcomputers.

HOLD1, FLUP, DONTFIX, FIXIT, and DSKSEL work together to read the disk drive selection, filename, and filetype and load them into the FCB.

Assuming that the proper wire connections have been made, the next step in Master is to send an 'R' on line to get the receiving micro's attention. Then the sending micro listens for a reply from the receiving micro. This is repeated until the sending micro receives an 'r' in reply. Master then prints a string to the screen to tell the user that connection has been made.

To ensure synchronization prior to sending the FCB, Master sends a Transmit Symbol (TXSYM). We use the ASCII equivalent for a DC2 control code as our TXSYM, chosen based on our determination that DC2 is not used

frequently otherwise. Master then listens for a reply. As a buffer, this is repeated until the sending micro receives a 't' in reply. Before sending the FCB, an open file subroutine is called to insure that the file exists. If the file exists, the program continues. Otherwise, the session is aborted through a 'FNFOUND' subroutine. A 'QUIT' symbol, the ASCII Code for a DC4 control code, is sent online to tell the receiving micro that no file transfer will occur. Then a string is printed to the screen telling the user that no file was found and the program returns to CPM.

We use the B register to store the current checksum code, initializing it to zero (0) for reference. The HL register pair holds the address of the current memory (M) location for purposes of data manipulation. To send the FCB, we set the pointer in the HL pair at the starting memory location for the FCB (005C Hex). The next loop uses the current memory byte to perform the checksum operation and sends that byte on line until the current memory location holds a '0'. Once that '0' is sent on line, the loop is done, as the '0' denotes the end of that filename. The checksum code is a result of 'exclusive oring' the current data byte with the previous checksum code. The resulting checksum code is stored back in the B register. Use of a checksum ensures accurate data transmission.

After the end of filename has been sent, the sending micro waits for an 'r' indicating that the receiving micro received the end of file '0' signal. The checksum is then sent online. We save the checksum for possible retransmission, then clear the accumulator before listening for acknowledgement. If a 'b' is received, the checksums didn't match, so the FCB is resent using RSNDFCB. First, the checksum is recalled from the stack and moved to the

accumulator. Then we offset the checksum by adding three (3) for use in synchronizing the two micros, and send the result online. Next is an indefinite wait loop that is left only when the reply matches RXACK ('r'). Following is a similar loop listening for a TXACK ('t'). When synchronization is set, the program jumps back to the subroutine TXFCB1 and resends the FCB. If a 'g' is received in reply, the transmitting micro proceeds to a wait loop for the receiving micro to catch up.

In the wait loop, the program checks for an input as many as 2000 times. If no input is received, the user is returned to CPM. When an input is received, it is compared to 'QUIT', a DC4 in ASCII Code. If a match is made, it means that the receiving disk already has a file of the same name and the program jumps to 'GOCPM1'. Here, a string is printed to the screen telling the user that the receiver already has a file of the same name and the user is returned to CPM. If the reply wasn't a 'QUIT', it is compared to a 'GCCN' or continue symbol. If the input matches neither of the two, the wait loop is repeated; otherwise, a string is printed to the user screen that the file is being transmitted. Next, the program calls a read sequential subroutine to get the first(next) 128-byte block of data.

Prior to sending each 128-byte block, a 'CHECK' subroutine is called see if the sending micro is ready to transmit. 'CHECK' holds the program until the micro is 'transmit ready'. Then, for synchronization, a TXSYM is sent online. A listen loop follows, where the program checks for a TXACK or a disk full symbol (DSKFUL), which is a 'd'. If TXACK is received, data can be sent. If DSKFUL is received, it means that the receiving micro has no more disk storage space and a full disk (FULDISK) abort subroutine is called. The FULDISK subroutine

sends a DONE symbol, a 'Z', online to acknowledge the 'DSKFUL' symbol. Then a string is printed to the screen telling the user that the receiver's disk is full, and the subroutine 'GOCPM' is called. First, a '0' is sent online to clear the output buffer of the sending micro and the input buffer of the receiving micro. Then the program returns to CPM. We found it necessary to send the '0' in order to prevent premature synchronization by the Slave micro. When we allowed the micro to return to CPM without this step, the Slave micro acted on whatever was left in the Master output buffer. This synchronization sequence is repeated until a match is made on one of the two expected inputs.

To separate the 128-byte frame from our control commands, MASTER now sends a Real Data (RLDTA) symbol, 0CB Hex, to the receiving micro. MASTER then listens for an echo from SLAVE before continuing with file transmission. Again, this was necessary for synchronization between different types of computers.

When an echo is received, we set the H,L register pair pointer to the first location in the Direct Memory Address (DMA) buffer, which is 80 Hex. The DMA is 80 Hex, or 128 bytes of memory, and is the default storage location for data read to or from files by CPM. The checksum in register B is reset to 0 for each 128-byte block. Now a checksum is performed in the same manner as it was for the FCB, and the current byte is moved to the accumulator to be sent online. Then the H,L pointer is moved to the address of the next data byte in the DMA. This is repeated until the 128th byte is sent and the H,L pointer is incremented to 100Hex. When the last byte of the data block has been sent, the checksum is moved to the accumulator and sent online.

Next, we have another listen loop to allow the receiving micro to catch up. The program checks for input until one is received. Once an input is received, it is compared against the 'Bad' and 'Good' symbols. If it is 'Bad', the program jumps to a 'RESEND' subroutine. In 'RESEND', a 'b' is printed to the user's CRT telling them that the block checksum was bad and that same 128-byte block is to be resent. Then the block is sent again. If it is a 'Good', the program jumps to a subroutine to send the next 128-byte block, 'RDSQRPT'. Here, a '*' is printed to the user's CRT telling them that the block was successfully sent. Then the program jumps to 'RDSEQ' to read the next block of data to be sent. If there is no more data in the file, a TXSYM is sent online. The program then listens for a TXACK until one is received. Then 'EOFIL1' is called. First, a QUIT symbol is sent online. Then the program listens for an echoing QUIT symbol, repeating until the echo is received. Then a string is printed to the screen telling the user the file transfer is complete and 'CLOSIT' is called. A DONE symbol is sent online, the transmitting micro listens for an echoing DONE in reply. Once the DONE echo is received, the program returns to CPM after sending a '0' online to clear the buffers. The listen sequence is repeated up to 26 times. If no match is made on 'Bad' or 'Good', we assume a problem and send the same 128-byte block again using the same procedures.

The 'POUT1' subroutine includes the ability for the sending user to abort file transfer. At any time during the program the user can enter a 'Control C' (ctrl C) from the keyboard to abort. 'POUT1' looks for this input every time it is called and, if 'ctrl C' is found, jumps to a 'STOPS' subroutine. The subroutine sends a CTRLC symbol online, then clears the accumulator and listens for a CTRLC

echo from the receiving micro. This is repeated until an echo is received. The output buffer is then cleared and the program returns the user to CPM.

There is one subroutine, 'OUTPUT', that is not used actively in the program. OUTPUT is left in the program code for debugging purposes in future revisions. This subroutine prints whatever is in the accumulator to the screen. Thus, the programmer can compare what is there against what was expected. We used this subroutine heavily in writing the program code.

The parallel part of the program that coordinates with the MASTER section is SLAVE. In order for MASTER to operate correctly on the initiation end of data transfer, the receiving end must have a working copy of MICROLAN on his disk. The following documentation will be a description of how SLAVE works in conjunction with MASTER.

In order for the receive portion (SLAVE) of the program to be initiated, the receiving operator must initialize his copy of MICROLAN. As previously stated, the program is executed by typing the word "MICROLAN". The operator will then be prompted to identify which disk drive he desires to work from, (A,B,C, or default), and then be prompted for an "R" to initiate the execution.

The program begins by listening for an attention signal, which is an 'R' (ATTN) from the transmitting micro. This is used by the MASTER to see if someone is out there ready to accept data transfer. SLAVE continues to listen until an ATTN is received. Once it is received, a message string is printed to the screen to let the operator know that a connection was made. SLAVE then sends an 'r' (RXACK) to the MASTER to acknowledge receipt of the ATTN.

The same procedure is essentially repeated, only with a few changes. SLAVE now listens for a 'DC2' (TXSYM) and

continues to listen until one is received. This is done for synchronization and acknowledgement that SLAVE is aware that data transfer is about to take place.

Before an acknowledgement signal is sent back to MASTER, a few operations will take place. This is done for synchronization.

The filename of the file that is being transferred is stored in a memory location known as the File Control Block, of FCB. The size is 32 spaces. The FCB is reset with zeroes to ensure that any previous data will not interfere. Once the FCB is reset, a 't' (TXACK) is sent to MASTER for acknowledgement that synchronization is set and SLAVE is ready for data reception.

The filename will be the first bit of information sent. Once SLAVE receives that first byte, it does a few comparisons before it writes it to memory.

First it checks for a 'DC4' (QUIT). If it receives one of these, it prints a message to the screen stating that no file transfer has taken place and then jumps out of the program (back to CPM). If the data was not a QUIT, then it is compared to a zero. A zero means that the filename has been completely sent and the program continues. If it was not a zero, then the comparison is against a TXSYM. This is done to ensure that the data was valid. A few TXSYM's may have been sent over from MASTER after synchronization was established on the SLAVE end. This procedure is a safeguard against reading those extra TXSYM's as data. If one does get through, the program loops itself until valid data is received.

Once the filename data is received, it is put into the FCB memory location and then printed to the screen. This allows the operator to see which file is being sent. A checksum is calculated (see MASTER for explanation of method) throughout the reception of the filename for use

later as a verification that the correct filename was received.

After the filename is received, an RXACK is transmitted to the MASTER to acknowledge that the filename has been received and SLAVE is awaiting the checksum calculated by MASTER. When this data is received, it is compared to the checksum calculated by SLAVE. If they are not the same, three (3) is added to the value sent by MASTER to announce that the checksums did not match. This means that the filename sent was not the same as the filename received. SLAVE then awaits for a re-transmission of the checksum + 3 sent previously. Once this is received, SLAVE acknowledges with a RXACK (which ensures synchronization), returns to reset the FCB and starts all over again, listening for a re-transmission of the filename.

If the checksums do match, then the program continues by sending a 'g' (GOOD). This verifies receipt of a good checksum. The subroutine OPNFILE then checks the directory to see if a file already exists by that filename which was previously transmitted. If one does exist, a QUIT is sent to MASTER advising that micro that a file already exists by that filename. A string is then printed to the screen telling the operator of the duplication of filenames, followed by the program jumping to CPM, terminating this session of SLAVE. If the filename did not previously exist in the directory, then a new file is created.

We are now ready to receive the data in 128-byte blocks. The Direct Memory Address (DMA) is a dedicated block of memory, 128 bytes long, used for this purpose. A synchronization check is done first and then a TXACK is sent to MASTER when SLAVE is ready to receive data. To separate the received data from MICROLAN's command and synchronization bytes, SLAVE now looks for a RLDTA symbol from MASTER. While SLAVE is looking for RLDTA, it also

performs one other check. If a QUIT was sent, that signals the end of transmission and the file is closed. Once SLAVE sees a RLDTA symbol, it echoes back to MASTER, then proceeds to look for data. First, SLAVE checks to see that RLDTA is not still being sent by MASTER (only on the first byte of the 128-byte block). Once SLAVE is sure that the data block is being sent, it enters the data receive loop. The data byte is moved into memory and a checksum calculated for each run through this loop. This procedure continues until the counter, initialized with the size of the DMA, has reached zero. This indicates that 128 bytes of data has been sent. MASTER sends its checksum and SLAVE compares it with its own. If it does not match, SLAVE sends a 'b' (BAD) to MASTER indicating that it must re-transmit the same 128 bytes. If the checksums do agree, then the 128 bytes are written to the disk and an asterisk is printed to the screen telling the operator that 128 bytes of data have been successfully transferred. The program then returns to repeat this process until a QUIT is received.

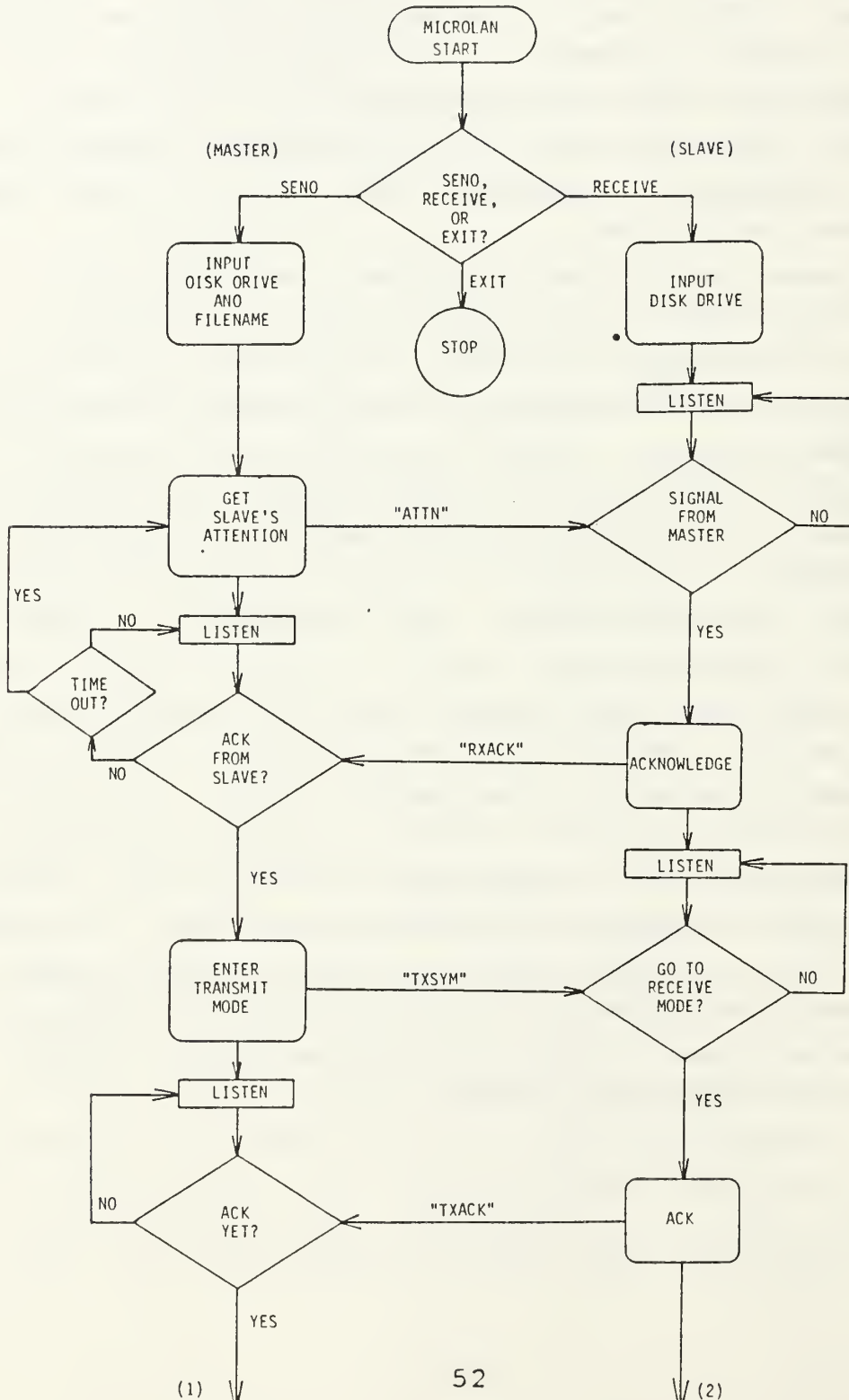
When a QUIT is received, SLAVE acknowledges by sending a QUIT back and then closes the file. A string is printed on the screen indicating to the operator that file transmission is complete. SLAVE then waits for a 'Z' (DONE) from MASTER which ensures that the session is complete. A DONE is transmitted back which completes the hand-shaking process and then SLAVE jumps to CPM. The SLAVE program has been terminated and the micro is ready for any command. If the operator wishes to receive another file, he must reinitiate the MICROLAN program.

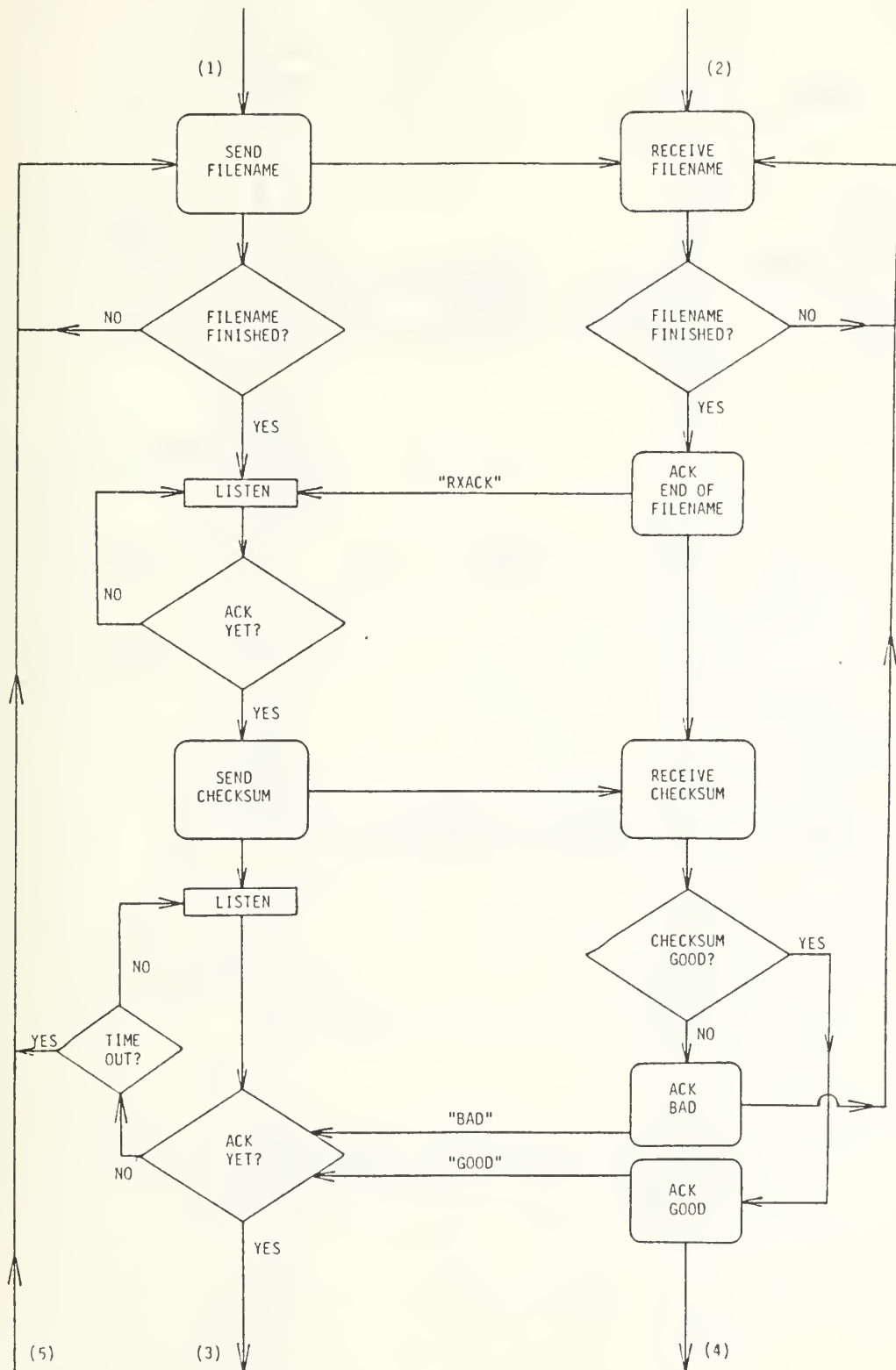
There are two safety factors that are included in the SLAVE program that were not previously mentioned. The first one concerns the occurrence of a full disk on the part of the receiving micro. Each time the program

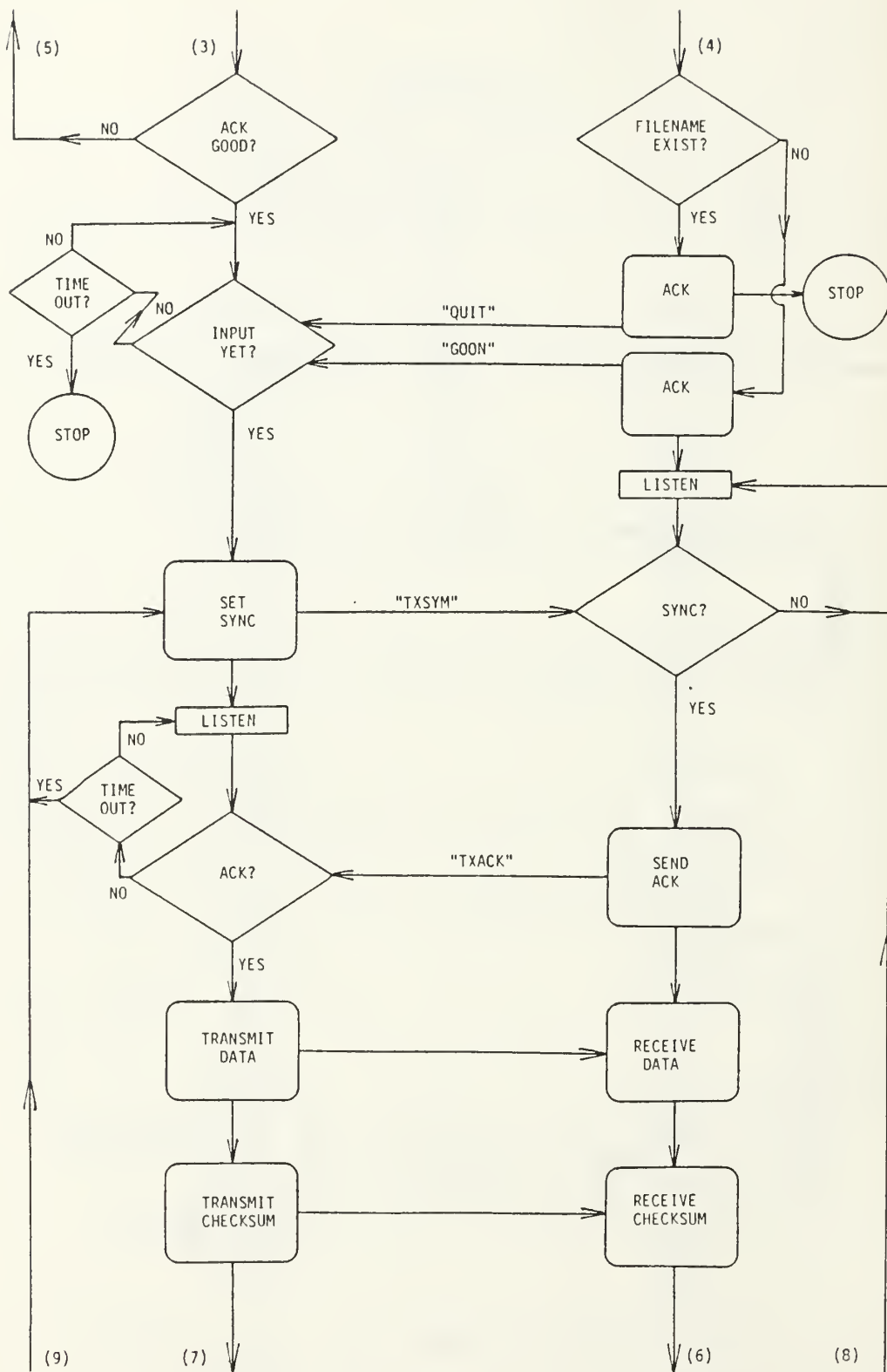
writes a 128-byte block of data to the disk, it checks to see if the disk is full. In the event of a full disk, SLAVE sends a 'd' (DSKFUL) to MASTER expressing that there is no more room on the disk and cannot receive any more data. SLAVE then awaits confirmation from MASTER that it has received the DSKFUL. Confirmation is acknowledged by the receipt of a DONE, which completes the "handshaking". A string is printed to the screen letting the operator know that he received an incomplete file due to a full disk. SLAVE then goes to CPM.

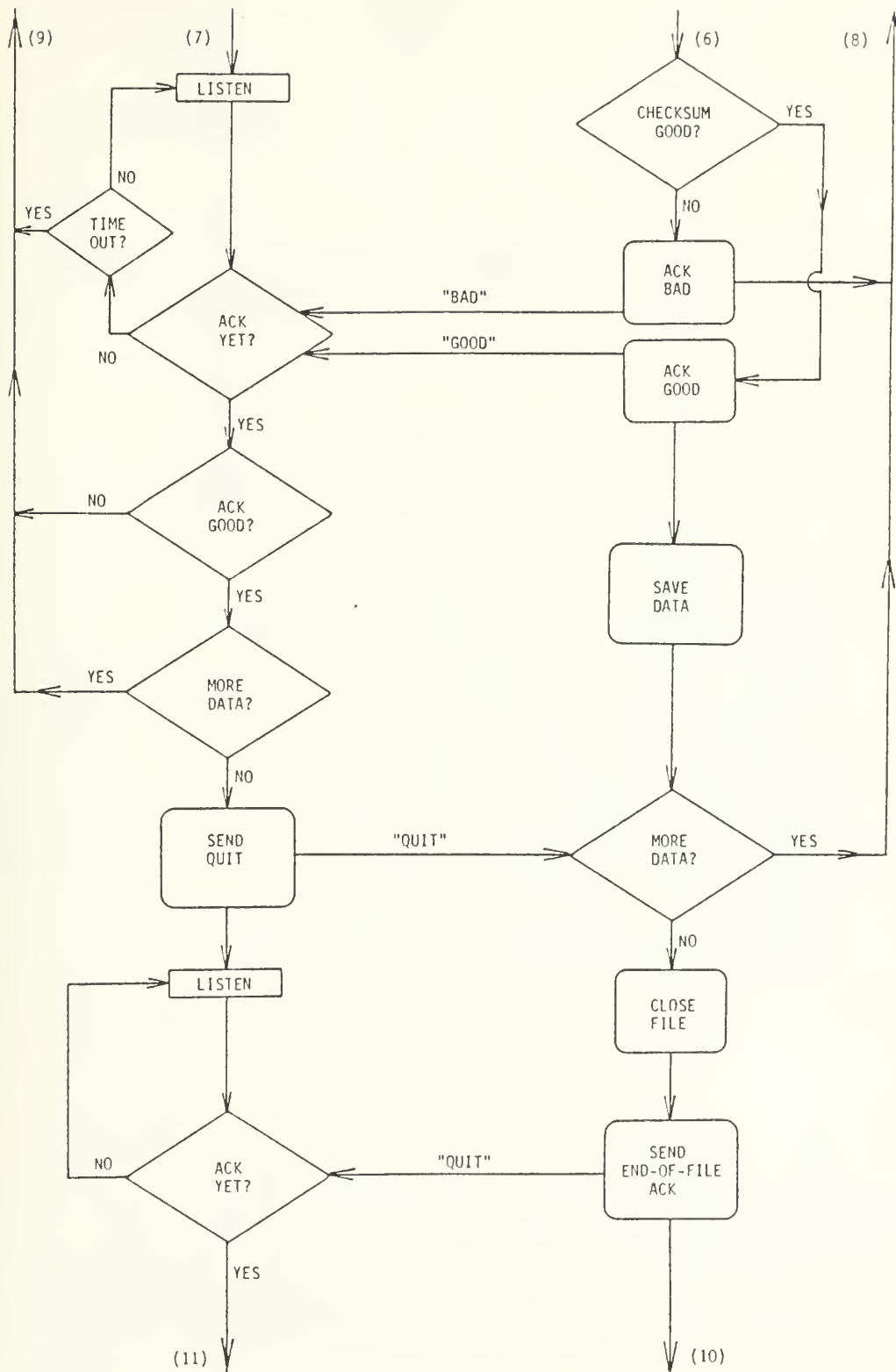
The other safety factor handles the occurrence of the transmitting micro aborting file transfer. To abort file transfer, the operator of the transmitting file uses a "control C". SLAVE listens for this "control C" throughout the entire program. Every time data is received from MASTER, it is checked for the abort signal. This allows for the option of the operator at the transmitting micro to stop data transfer at any time. If this happens, the program goes to a subroutine which sends a "control C" back to MASTER in acknowledgement and then prints a string to the screen. This tells the operator that file transfer has been aborted and that no file will exist under the filename that was passed. The program then jumps to CPM. Our logic was based on "whole file or no file". We felt that having an empty file would be an unmistakable indicator that the file transfer was incomplete and that retransmission was necessary. If the operator wishes to retain a partial file, a minor change to the program would be needed. The file would have to be closed before the program jumped to CPM by invoking the subroutine "CLSFILE" first.

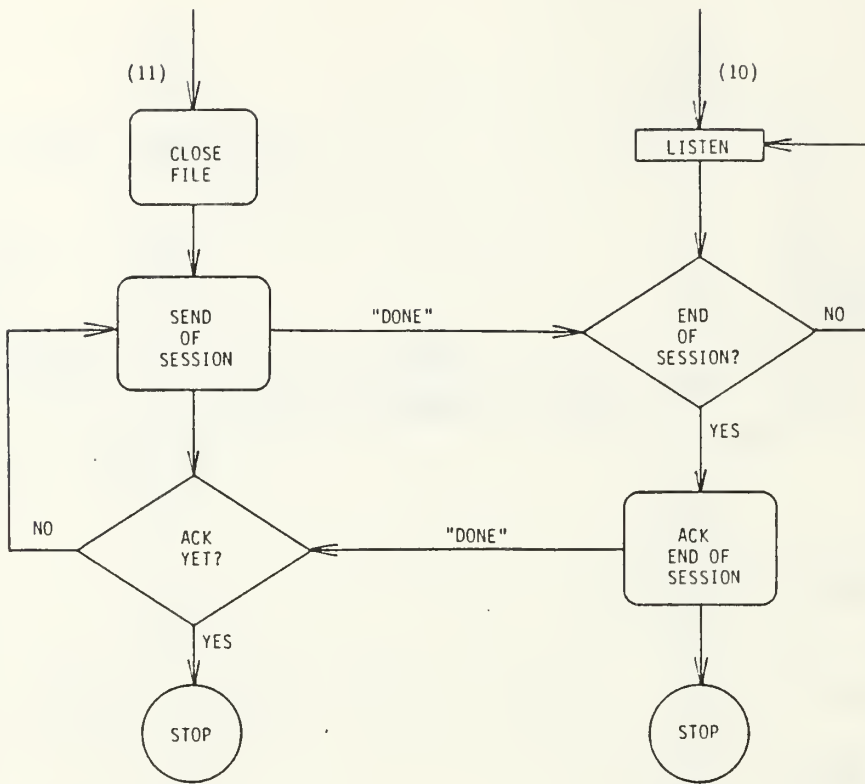
APPENDIX B FLOW DIAGRAM











APPENDIX C

CP/M - 80

```

TRUE      EQU 0FFFFH
NORTHSTAR EQU TRUE
APPLE     EQU NOT TRUE
ORG 0100H

INIT      LXI SP,0A000H      ;MCVE STACK PTR TO SAFE LOC
          CALL CRLF
          MVI C,09H          ;PFIINT STRING TO SCREEN
          LXI D,RIGHTS       ;CCPYRIGHTS AND NAMES OF AUTHORS
          CALL BDOS
          CALL CRLF
          CALL CRLF
          MVI C,09H          ;PFIINT STRING TO SCREEN
          LXI D,WELCUM       ;WELCOME MSG
          CALL BDOS
          CALL CRLF
          CALL CRLF
          MVI C,09H          ;PFIINT STRING TO SCREEN
          LXI D,INSTRC       ;SEND, RECEIVE, OR QUIT?
          CALL BDOS
          CALL CRLF
          CALL CRLF
```

```

HOLDING MVI C,06H
MVI E,0FFH
CALL BDOS
ANI 0DFH
CPI 53H
JZ MASTER
CPI 52H
JZ SLAVE
CPI 58H
JZ CPM
JMP HOLDING
MASTER MVI C,09H
IXI D,ENTER
CALL BDOS
CALL CRLF
CALL CRLF
FILLUP LXI H,FCB
MVI M,00H
INX H
MVI C,0BH
FILLUP1 MVI M,20H
INX H
DCR C
JNZ FILLUP1
;CHECK FOR CONSOLE INPUT
;LOOKING FOR INPUT
;ENSURE LETTER IS A CAPITAL
;IS IT AN 'S'?
;IF SO, START FILE TRANSFER
;IS IT AN 'R'?
;IF SO, PREPARE TO RECEIVE FILE
;IS IT AN 'X'?
;IF SO, GO TO CPM
;REPEAT UNTIL INPUT FOUND
;PRINT STRING TO SCREEN
;ENTER FILENAME
;ADDRESS OF FCB
;11 SPACES
;FILL MEMORY ADDRESS WITH SPACES
;MOVE PTR TO NEXT ADDRESS
;DECREMENT COUNTER
;REPEAT UNTIL DONE

```

FILLUP2	MVI C,13H		;TCTAL OF 20 SPACES
	MVI M,00H		;FILL REST OF ADDRESS WITH 0'S
	INX H		;MCVE PTR TO NEXT ADDRESS
	DCR C		;DECREMENT COUNTER
	JNZ FILLUP2		;REPEAT UNTIL DONE
HOLD1	MVI C,0AH		;READ CONSOLE BUFFER
	LXI D,CONBUF		;ADDRESS OF FIRST LETTER OF FILENAME
	CALL BDOS		
	LXI H,CONBUF		;ADDRESS OF CONSOLE BUFFER
	LXI D,FCB+1		;FCB ADDRESS
	INX H		
	MOV B,M		;STORE COUNT IN B REGISTER
	MOV A,M		;MCVE COUNT TO ACCUMULATOR
	ORA A		;IS THERE AN INPUT?
	JZ ERROR		;TRY AGAIN
	INX H		
FLUP	MOV A,M		
	CPI 3AH		;IS CHARACTER A ':'?
	JZ DSKSEL		;IF SO, GO TO DISK SELECT
	CPI 2EH		;IS IT A '.'?
	JZ FIXIT		;IF SO, SKIP TO FILETYPE
	CPI 40H		;CHECK FOR LETTER
	JC DONTFIX		;SKIP NEXT STEP IF NOT LETTER
	ANI 0DFH		;ENSURE LETTER IS A CAPITAL

```

DONTFIX STAX D      ;STORE CHARACTER IN FCB
INX D
INX H
DCR B

JNZ FLUP
LISN1 MVI A,ATTN
      CALL POUT1
      MVI C,03H
      CALL PIN
      CPI RXACK
      JZ XMIT
      DCR C
      JNZ LISN
      JMP LISN1
      CALL SWAP
      CALL CRLF
      CALL CRLF
      MVI C,09H
      LXI D,RXING1
      CALL BDOS
      CALL CRLF
      CALL SWAP
      XMIT1 MVI A,TXSYM
            CALL POUT1
            ;DC2 SYMBOL FOR SYNC AT START
            ;OF 128 BYTE BLOCK

```



```

MVI C,08AH
LITTLT CALL PIN
CPI TXACK
JZ TXFCR
DCR C
JNZ LITTLT
JMP XMIT1
TXFCB CALL CRLF
CALL OPENIT
TXFCB1 MVI B,00H
LXI H,FCB
FCBLUP MOV A,B
INX H
XRA M
MOV B,A
MOV A,M
CALL POUT1
CPI 0H
JZ FCBCK
JMP FCBLUP
FCBCK MVI C,20H
FCBCK1 CALL PIN
CPI RXACK
JNZ FCBCK1

;LISTEN 138 TIMES

;WAS 't' RECEIVED?
;IF SO, XMIT FILE CTRL BLK
;OTHERWISE KEEP LISTENING
;UNTIL CTR IS ZERO,
;TEEN SEND DC2 SYNC AGAIN

;SEE IF FILE EXISTS. IF SO, OPEN IT
;INITIALIZE CHECKSUM REGISTER
;SET PTR TO 1ST LETTER IN FILENAME
;PERFORM CHECKSUM OPERATION
; (MOVE PTR TO NEXT BYTE)
;BY XORING CURRENT BYTE
;WITH B REGISTER
;PUT CURRENT BYTE IN ACCUM
;SEND CURRENT BYTE
;CHECK FOR END OF FILENAME
;IF END, GO TO CHECKSUM LOOP
;IF NOT, REPEAT FCB LOOP
;LCOP 32 TIMES
;FCR SYNC WITH SLAVE
;IS IT AN 't'?
;IF NOT, LISTEN AGAIN. IF SO,

```

MOV A,B	;PUT CHECKSUM IN ACCUM
CALL POUT	;SEND CHECKSUM
PUSH B	;SAVE CHECKSUM
MVI A,0H	;CLEAR ACCUM
MVI B,80H	;LISTEN 100 TIMES
FCBTFMOT CALL PIN	;READ MAIL
CPI BAD	;DID IT CHECK BAD?
JZ RSNDFCB	;IF SO, SEND FCB AGAIN
CPI GOOD	;DID IT CHECK GOOD?
JZ WAITFIL	;IF SO, GO TO NEXT ROUTINE
DCR B	;IF NOT, DECREMENT CTR, AND
JNZ FCBTFMOT	;IF NOT 0, LISTEN AGAIN
POP B	;CLEAR STACK
DCR C	;IF SO, DECREMENT C
JNZ FCBCK1	;AND REPEAT UNTIL C=0
JMP TXFCB	;IF 0, ASSUME PROBLEM AND SEND AGAIN
WAITFIL POP B	;CLEAR STACK
WAIT2 LXI B,07FFH	;CCUNT LOOP APPX 2K
WAIT1 CALL STATIN1	;ANY 'MAIL'?
JZ WAIT1	;IF NOT, CHECK AGAIN
DCX B	;IF SO, DECREMENT CTR
MOV A,B	
CRA C	
JZ GOCPM	;AND, IF 0, QUIT

CALL PIN	;OTHERWISE READ 'MAIL'
CPI QUIT	;DCES RXING MICRO ALREADY HAVE FILE?
JZ GOCPM1	;IF SO, GO TO CPM
CPI GOON	;IS IT THE GO ON SIGNAL 'G'
JNZ WAIT2	;IF NOT, LISTEN AGAIN. ALLOW RXING
	;MICRO TO CATCH UP
CALL SWAP	
CALL CRLF	
CALL SWAP	
TXDATA CALL SWAP	;SEND THE FILE
MVI C,09H	;PFINT STRING TO SCREEN
LXI D,TXING1	;SAYS FILE BEING SENT
CALL BDOS	
CALL CRLF	
CALL SWAP	
RDSEQ CALL READSEQ	;READ FIRST 128 BYTE BLOCK
SEND CALL CHECK	;AND SEND TO RXING MICRO
MVI A,TXSYM	;DC2 SYMBOL FOR SYNC AT START OF DATA
CALL POUT1	
MVI C,0FH	;LISTEN 15 TIMES
LITLET2 CALL PIN	
CPI TXACK	;IS IT A 't'
JZ SLUP2	;IF SO, READY TO SEND DATA
CPI DSKFUL	;IS RXING MICRO'S DISK FULL?

JZ FULDISK	; IF SO, QUIT
DCR C	; IF NOT, DECREMENT CTR
JNZ LITLET2	; LISTEN AGAIN, UNLESS CTR IS 0,
JMP SEND	; THEN TRY TO SYNC AGAIN
SLUP2 MVI A,RLDTA	; OCBH MEANS TIME FOR DATA
CALL POUT1	
LXI B,07FFH	; WAIT LOOP APPX 2K
SLUP3 CALL PIN	
CPI RLDTA	; IS IT ECHO?
JZ SLUP1	; IF SO, SEND DATA
DCX B	; DECREMENT COUNTER
MOV A,B	
CRA C	
JNZ SLUP3	; REPEAT UNTIL ZERO
JMP SLUP2	; THEN SEND AGAIN
SLUP1 LXI H,DMA	; PCINTER TO 1ST INFO BYTE
MVI B,00H	; INITIALIZE CHECKSUM
SLOOP MOV A,B	; PERFORM CHECKSUM
XRA M	; XCR DATA WITH B REGISTER
MOV B,A	
MOV A,M	; PUT BYTE IN ACCUMULATOR
CALL POUT	; DATA IS TRANSFERRED
INX H	; MCVE PTR TO NEXT BYTE
MOV A,H	; MCVE H REG TO ACCUM

CRC	CPI ENDMA	;END DMA, CHECK FOR LAST BYTE
	JNZ SLOOP	;IF NOT, SEND NEXT BYTE. OTHERWISE
	MOV A,B	;PUT CHECKSUM IN ACCUMULATOR
	CALL POUT	;AND SEND TO RXING MICRO
CRCTMOT	MVI B,01AH	;LISTEN 26 TIMES
CRCT1	CALL STATIN1	;CHECK INPUT BUFFER
	JZ CRCT1	;IF NOTHING, TRY AGAIN
	CALL PIN	;READ MAIL
	CPI BAD	;IS CHECK BAD?
	JZ RESEND	;IF SO, SEND BLOCK AGAIN
	CPI GOOD	;IS CHECK GOOD?
	JZ RDSQRPT	;IF SO, READ NEXT BLOCK
	LCR B	;DECREMENT COUNTER
	JNZ CRCT1	;IF NOT TIMED OUT,LISTEN AGAIN
	JMP SEND	;IF TIMED OUT, ASSUME PROBLEM, AND
		;SEND BLOCK AGAIN
DSKSEL	CALL SWAP	
	IXI H,CONBUF+2	;ADDRESS OF DISK SEL ENTRY
	MOV A,M	;PUT DISK SEL IN ACCUM
	ANI ODFH	;ENSURE LETTER IS CAPITAL
	CPI 'A'	;IS LETTER AN 'A'?
	JZ ADISK	;IF SO, SET FOR A DRIVE.
	CPI 'B'	;IS LETTER A 'B'?
	JZ BDISK	;IF SO, SET FOR B DRIVE.

```

ADISK      JNZ DSKSEL1
           LXI H,FCB
           MVI M,01H
           JMP DSKSEL1
BDISK      LXI H,FCB
           MVI M,02H
           CALL SWAP
           LXI H,CONBUF+4
           LXI D,FCB+1
           JMP FLUP
           CALL SWAP
RESEND     CALL SWAP
           MVI C,CONOUT
           MVI E,BAD
           CALL BDOS
           CALL SWAP
           JMP SEND
RSNDFCB   POP B
           MOV A,B
           ADI 3
           CALL POUT1
           CALL PIN
           CPI PYACK
           JNZ RSNDF
           RSNDFC1 CALL PIN

           ; IF NEITHER, RETURN TO FILENAME LOOP.
           ; SET PTR TO DRIVE BYTE.
           ; SET FCB FOR A DRIVE.
           ; RETURN TO FILENAME LOOP.
           ; SET PTR TO DRIVE BYTE.
           ; SET FCB FOR B DRIVE.

           ; MOVE BUFFER POINTER TO FILENAME.
           ; FCB FILENAME ADDRESS.
           ; RETURN TO FILENAME LOOP.

           ; PRINT TO SCREEN
           ; A 'b' IF CHECKSUM WAS BAD

           ; AND SEND BLOCK AGAIN
           ; RECALL CHECKSUM
           ; PUT CHECKSUM IN ACCUM
           ; ADD 3 TO OFFSET
           ; SEND BYTE

           ; IS IT AN 'r'?
           ; IF NOT, LISTEN AGAIN
           ; READ MAILBOX

```

CPI TXACK	;SYNC WITH RXING MICRO
JNZ RSNDFC1	;REPEAT UNTIL TXACK RECEIVED
JMP TXFCB1	;IF SO, RESEND FCB
LXI D,FCB+9	;MCVE POINTER TO FILETYPE AREA
INX H	;MCVE PTR TO FIRST LETTER OF FILTYP
JMP FLUP	
MVI C,09H	;PFINT STRING TO SCREEN
LXI D,ERNMSG	;ERROR MESSAGE
CALL BDOS	
CALL CRLF	
JMP HOLD1	;LCOK FOR INPUT AGAIN
PUSH PSW	;SEND THE DATA
	;FIRST, SAVE THE CURRENT BYTE
MVI C,06H	;CHECK FOR CONSOLE INPUT
MVI E,0FFH	;LCOKING FOR INPUT
CALL BDOS	
CPI CTRLC	;IS THERE A CONTROL C?
JZ STOPS	;IF SO, ABORT; OTHERWISE,
CALL CHECK	;PERFORM CHECK
POP PSW	;AND RECALL BYTE
IF NORTHSTAR	;IF MICRO IS NORTHSTAR
OUT DATA	;THEN XMIT BYTE
ENDIF	
IF APPLE	;IF MICRO IS APPLE

STA DATA1	;XMIT BYTE
ENDIF	
RET	
OPENIT	;OPEN FILE CODE
LXI D,FCB	;FILE CTRL BLOCK ADDRESS IN DE REG PR
CALL BDOS	
CPI OFFH	;FF = FILE NOT FOUND
JZ FNFOUND	;IF FILE NOT FOUND
RET	;CTHERWISE, RET TO TX DATA
	;CLOSE FILE CODE
CLOSLIT	;FILE CTRL BLOCK ADDRESS IN DE REG PR
LXI D,FCB	
CALL BDOS	
CLOSLIT1	;END OF SESSION MSG 'Z'
MVI A,DONE	;SEND TO RXING MICRO
CALL POUT1	;CLEAR ACCUM
MVI A,0H	;CHECK REPLY
CALL PIN	;DCES RXING MICRO AGREE?
CPI DONE	;IF NOT, REPEAT
JNZ CLOSLIT1	;IF SO, GO TO CPM
JMP GOCPM	;READ SEQUENTIAL CODE
READSEQ	;FILE CTRL BLOCK ADDRESS IN DE REG PR
MVI C,14H	
LXI D,FCB	
CALL BDOS	
CPI 0	;0 MEANS SUCCESSFUL READ
JNZ EOFILE	;IF NOT 0, ASSUME FINISHED WITH FILE

```

RDSQRPT CALL SWAP
RET
MVI C,CONOUT
MVI E,02AH
CALL BDOS
CALL SWAP
JMP RDSEQ
FNFOUND MVI A,QUIT
CALL POUT1
MVI C,09H
LYI D,FNFDMSG
CALL BDOS
CALL CRLF
JMP GOCPM
EOFIL1 POP PSW
EOFIL2 MVI A,TSYX
CALL POUT1
MVI C,0FH
LITLET3 CALL STATIN1
JZ LITLET3
CALL PIN
CPI TXACK
JZ EOFIL1
DCR C
;PRINT TO SCREEN
;*' SO USER KNOWS BLK WAS SENT
;TC READ NEXT 128 BYTE BLK
;TELL RXING MICRO NO FILE FOUND
;PRINT STRING TO SCREEN
;FILE NOT FOUND MSG
;AND GO TO CPM
;CCRECT STACK POINTER
;DC2 SYMBOL FOR SYNC WITH RXING MICRO
;LISTEN 15 TIMES
;CHECK FOR MAIL
;IF NONE, CHECK AGAIN
;READ MAIL
;IS IT A 't'?
;IF SO, CONTINUE
;IF NOT, DECREMENT COUNTER

```



```

JNZ LITL3
JMP EOFIL2
EOFIL1 MVI A,QUIT
CALL POUT1
CALL PIN
CPI QUIT
JNZ EOFIL1
CALL CRLF
MVI C,09H
LXI D,EOFMSG
CALL BDOS
CALL CRLF
JMP CLOSIT
STOPS MVI A,CTRLC
IF NORTHSTAR
    OUT DATA
    ENDF
    IF APPLE
        STA DATA1
    ENDF
    MVI A,0H
    CALL PIN
    CPI CTRLC
    JNZ STOPS
;AND LISTEN AGAIN, UNLESS COUNTER IS
;0. THEN TRY AGAIN
;DC4 SYMBOL. TELLS RXING MICRO THAT
;THE FILE IS DONE
;LISTEN FOR REPLY
;DCES RXING MICRO ACKNOWLEDGE?
;IF NOT, TRY AGAIN
;PRINT STRING TO SCREEN
;IF SO, TELL USER FILE IS DONE
;AND CLOSE THE FILE
;SEND CTRLC TO RXING MICRO
;IF NORTHSTAR MICRO
;SEND CONTROL C
;IF APPLE MICRO
;SEND CONTROL C
;CLEAR ACCUM
;FROM RXING MICRO
;ACK FROM RXING MICRO
;REPEAT UNTIL ACK

```

```

POP PSW
JMP GOCPM

FULDISK MVI A,DONE
CALL POUT1
MVI C,09H
LXI D,FULMSG
CALL BDOS
CALL CRLF
JMP GOCPM

GOCPM MVI A,0H
IF NORTHSTAR
OUT DATA
ENDIF
IF APPLE
STA DATA1
ENDIF
CALL CRLF
CALL PIN
JMP CPM

GOCPM1 MVI C,09H
LXI D,HASFILE
CALL BDOS
JMP GOCPM

SLAVE MVI C,09H
;CORRECT STACK POINTER

;LETTER 'Z' TO ACKNOWLEDGE
;SEND BYTE
;PRINT STRING TO SCREEN
;SAYS RXER'S DISK FULL

;RESET THE ACCUMULATOR AND
;NCRTHSTAR MICRO
;SEND BYTE

;APPLE MICRO
;SEND BYTE

;AND GO TO CPM
;PRINT STRING TO SCREEN
;RXING MICRO HAS FILE ALREADY

;PRINT STRING TO SCREEN

```

LXI D, WCHDSK	;SELECT DISK DRIVE
CALL BDOS	
CALL CRLF	
CALL CRLF	
DRVSEL MVI C, 06H	;CHECK FOR CONSOLE INPUT
MVI E, 0FFH	;LOOKING FOR INPUT
CALL BDOS	
CPI 0DH	;IS IT A <CR>?
JZ CONT	;IF SO, ENTER RECEIVE MODE
ANI 0DFH	;ENSURE LETTER IS A CAPITAL
CPI 'A'	;IS IT AN 'A'?
JNZ DISKB	;SKIP TO B IF NOT 'A'
LXI H, FCB	;ADDRESS OF DISK DRIVE BYTE
MVI M, 01H	;SET BYTE TO A DISK DRIVE
JMP CONT	;THEN CONTINUE
CPI 'B'	;IS IT A 'B'?
JNZ DRVSEL	;IF NOT, CONTINUE IN RECEIVE MODE
LXI H, FCB	;ADDRESS OF DISK DRIVE BYTE
MVI M, 02H	;SET BYTE TO B DISK DRIVE
MVI C, 09H	;PRINT STRING TO SCREEN
CONT LXI D, RXMODE	;IN RECEIVE MODE
CALL BDOS	
CALL CRLF	
SLAVE1 MVI A, 00H	;RESET ACCUMULATOR

```

CALL PIN
CPI ATTH
JNZ SLAVE1

CALL SWAP
CALL CRLF
MVI C,09H
LXI D,RXING1
CALL BDOS
CALL SWAP
CALL CRLF
MVI A,RXACK
CALL POUT
LISTEN
CALL PIN1
CPI TXSYM
JNZ LISTEN

CALL CRLF
RXFCB
CALL SWAP
LXI H,FCB+1

MVI C,1EH
RSTFCB
MVI M,00H
INX H

;LISTENING FOR AN 'R'
; 'R'
;IF 'R' RX'D, CONTINUE. IF NOT
;LISTEN AGAIN

;PRINT STRING TO SCREEN
;CCNECTION MADE

; 'r'
;SEND A 'r' TO XMITING MICRO
;LISTENING FOR A 'DC2'
; 'IC2'
;IF 'DC2' RX'D, CONTINUE. IF NOT,
;LISTEN AGAIN

;ADDRESS OF FCB MEM LOC INTO
;H,L REGISTER PAIR
;CCOUNTER FOR FCB'S 31 SPACES
;FILL FCB WITH 0'S
;MOVE PTR TO NEXT MEMORY ADDRESS IN FCB

```

```

DCR C
MOV A,C
CPI 0
JNZ RSTFCB

CALL SWAP
MVI B,00H
IXI H,FCB+1

MVI A,TXACK
CALL POUT
CALL PIN1
CALL STATIN1
RST1
JZ RST1

CALL PIN1
RST2
CPI QUIT
JZ NOFILE
CPI 0H
JZ FCBCRC
CPI TMSYM
JZ RST2

MOV H,A
CALL OUTPUT

;DECREMENT COUNTER

;IF COUNTER = 0, CONT. IF NOT,
;PUT ANOTHER 0 IN FCB

;INITIALIZE CHECKSUM
;LCAD 2ND ADDRESS OF FCB IN
;H,L REGISTER PAIR
;'t'
;SEND 't' TO XMITING MICRC FOR SYNC
;CLEAR THE ACCUMULATOR
;CHECKING FOR INPUT

;FILE NAME DATA
;IS DATA A 'QUIT'?
;FILE DID NOT EXIST
;CHECK IF FILENAME COMPLETELY SENT
;IF FILENAME RX'D, GO TO CHECKSUM
;CHECK IF DATA IS VALID
;IF DATA IS NOT FILENAME,
;CHECK NEXT BYTE
;PUT FILENAME IN FCB
;PRINT FILENAME TO SCREEN

```



```

MOV A,B
XRA M
MOV B,A
INX H
JMP RST1

FCBCRC MVI A,RXACK
CALL POUT
FCBCRC1 CALL STATIN1
JZ FCBCRC1
CALL CRLF
CALL PIN
CMP B
JZ STRTFIL
ADI 3
MOV C,A
MVI A,BAD
CALL POUT
CLEAR CALL PIN
CMP C
JNZ CLEAR
MVI A,RXACK
CALL POUT
JMP EXFCB
STRTFIL MVI A,GOOD

;CALCULATE CHECKSUM
;MOVE PTR TO NEXT FCB ADDRESS
;'I'
;SYNC DATA WITH XMITING MICRO
;CHECKING FOR INPUT
;CHECKSUM DATA
;CCMPARE CHECKSUM
;CHECKSUM MATCHED
;ADD 3 TO THE CHECKSUM
;STORE IN REGISTER
;CHECKSUM DID NOT MATCH
;TELL XMITING MICRO
;XMITING MICRO STOPPED SENDING CHKSUM?
;IF NOT,LISTEN AGAIN
;SYNC WITH XMITING MICRO
;TRY AGAIN
;READY TO CHECK IF FILENAME ALREADY USED

```

RXD1	CALL POUT	
	CALL OPNFILE	;CHECK IF FILENAME EXISTS
	CALL MAKEFIL	;CREATE NEW FILE
	MVI B,00H	;INITIALIZE CHECKSUM
	LXI H,DMA	;LCAD ADDRESS OF DMA MEM LOC TO
		;H,L REGISTER PAIR
	MVI C,81H	;INITIALIZE COUNTER WITH SIZE OF DMA
RXDS	CALL STATIN1	;CHECKING FOR INPUT
	JZ RXDS	
RXD2	CALL PIN1	;SYNC WITH XMITING MICRO
	CPI TXXSYM	;CCMPARE WITH 'DC2'
	JNZ RXD2	
	MVI A, TXACK	; 't'
	CALL POUT	;IN SYNC WITH XMITING MICRO
RXDS1	CALL STATIN1	;CHECKING FOR INPUT
	JZ RXDS1	
RXYET	CALL PIN1	
	CPI RLDTA	;IS IT OCBH?
	JZ XXYET1	;IF SO, GO TO RECEIVE DATA
	CPI QUIT	;IS IT 'DC4' FOR QUIT?
	JZ CLSFILE	;IF SO, CLOSE FILE; OTHERWISE,
	JMP RXDS1	;LISTEN AGAIN
RXYET1	MVI A,FLDTA	;ACK REAL DATA COMING
	CALL POUT	

RXYET2	MVI A,00H		
	CALL STATIN1		;CLEAR ACCUM
RXYET3	JZ RXYET2		;CHECK FOR INPUT
	CALL PIN		
	CPI RLDTA		;READ DATA
	JZ RXYET3		;IS IT STILL RLDTA?
RXD3	DCR C		;DECREMENT COUNTER
	JZ RXCRC		;CHECKSUM RX'D
	MOV M,A		;PUT THE DATA IN MEMORY
	MOV A,B		
	XRA M		;CALCULATE CHECKSUM
	MOV B,A		
	INX H		;MCVE PTR TO NEXT DMA ADDRESS
RXD4	CALL STATIN1		;CHECK FOR INPUT
	JZ RXD4		;LCOP UNTIL INPUT
	CALL PIN		
	JMP RXD3		
RXCRC	MOV A,A		;ENSURE B IS COMPARED TO A
	CMP B		;CCMPARE WITH CHECKSUM
	JZ WRITFIL		;128 BYTE BLOCK SENT
	MVI A,BAD		;CHECKSUM DID NOT MATCH
	CALL POUT		;NCTIFY XMITING MICRO
	JMP RXD2		;SEND 128 BYTE BLOCK AGAIN
POUT	PUSH PSY		;SAVE THE DATA

```

CALL CHECK
POP PSW
IF NORTHSTAR
    OUT DATA
ENDIF
IF APPLE
    SIA DATA1
ENDIF
RET
PIN1 IF NORTHSTAR
    IN DATA
ENDIF
IF APPLE
    LDA DATA1
ENDIF
CPI CTRLC
JZ ABORT
RET
WRITEFIL MVI A,GOOD
CALL POUT
CALL WRITSEQ
CALL SWAP
MVI C,CONOUT
MVI E,02AH
;RETURN THE DATA
;MICRO IS NORTHSTAR
;SENDS DATA ACROSS THE LINE
;MICRO IS APPLE
;SENDS DATA ACROSS THE LINE
;MICRO IS NORTHSTAR
;MICRO IS APPLE
;DID XMITING MICRO ABORT?
;IF SO, ABORT
;XMIT THAT THE CHECKSUM IS CORRECT
;START WRITING FILE TO DISK
;SAVE REGISTERS
;PRINT TO SCREEN
; ** TO PRINT TO SCREEN

```

```

CALL BDOS
CALL SWAP
JMP RXD1
CPNFILE MVI C,0FH
LXI D,FCB
CALL BDOS
CPI 0FFH
JNZ FILFND
RET

CLSFILE MVI A,QUIT
CALL POUT
MVI C,10H
LXI D,FCB
CALL BDOS
CALL CRLF
LXI D,EOPMSG
MVI C,09H
CALL BDOS
CALL CRLF
MVI A,0H
CLSFILE1 CALJ PIN1
CPI DONE
JNZ CLSFIL1
MVI A,DONE

;RETURN REGISTERS

;OPEN FILE CODE
;FCB ADDRESS IN D,E RGSTR PAIR

;FF = FILE NOT FOUND
;FILE EXISTS

;'IC4'
;AGREE END OF FILE
;CLOSE FILE CODE
;FCB ADDRESS IN D,E RGSTR PAIR

;FILE TRANSMISSION COMPLETED
;PRINT STRING TO SCREEN

;CLEAR THE ACCUMULATOR
;LOOKING FOR END OF SESSION MSG
;'Z' = END OF SESSION

;END OF SESSION MESSAGE

```



```

CALL POUT
JMP CPM

MAKEFIL MVI C,16H
LXI D,FCB
CALL BDOS
MVI A,GOON
CALL POUT
RET

WRITSEQ MVI C,15H
LXI D,FCB
CALL BDOS
ORA A
JNZ FULLDSK
RET

FILFND MVI A,QUIT
CALL POUT
MVI C,09H
LXI D,FNDMSG
CALL BDOS
CALI CRLF
JMP CPM

NOFILE MVI C,09H
LXI D,NOMSG
CALI BDOS

;CCNFIRM RECEPTION OF E-O-SESSION MSG
;MAKE NEW FILE CODE
;FCB ADDRESS IN D,E RGSTR PAIR
;CCNTINUE MESSAGE
;RETURN TO RX FIRST 128 BYTE BLOCK
;WRITE THE FILE TO THE DISK
;FCB IN D,E RGSTR PAIR
;CHECK IF DISK IS FULL
;IF SO, JUMP TO FULLDSK
;TELL XMITING MICRO, FILE FOUND
;PRINT STRING TO SCREEN
;FILE ALREADY EXISTS. GO TO CPM
;PRINT STRING TO SCREEN
;NC FILE TRANSFER

```

```

CALL CRLF
JMP CPM

ABORT CALL CRLF
      MVI A,CTRLC
      CALL POUT
      MVI C,09H
      LXI D,ABRTMSG
      CALL BDOS
      CALL CRLF
      JMP GOCPM

FULLDSK MVI A,DSKFUL
        CALL POUT
        CALL PIN
        CPI DONE
        JNZ FULLDSK
        MVI C,09H
        LXI D,FULLMSG
        CALL BDOS
        CALL CRLF
        JMP CPM

CHECK  CALL STATIN2
      JZ CHECK
      RET

STATIN1 IF NORTHSTAR
      ; MICRO IS NORTHSTAR

      ;CHECK STATUS BYTE
      ;CCNTINUE UNTIL TXRDY IS SET

      ;PRINT TO SCREEN
      ;FILE TRANSFER INCOMPLETE, DISK FULL

      ;GC TO CPM
      ;
      ;TELL XMITTING MICRO DISK FULL
      ;AWAITING CONFIRMATION

      ;SEND XMITTING MICRO ABORT ACK
      ;PRINT STRING TO SCREEN
      ;XMITTING MICRO ABORTED

```

```

IN STATUS
ANI RXRDY
ENDIF
IF APPLE
;MICRO IS APPLE
LDA STATUS1
ANI RXRDY1
ENDIF
RET

STATIN2 IF NORTHSTAR
IN STATUS
ANI TXRDY
ENDIF
IF APPLE
;MICRO IS APPLE
LDA STATUS1
ANI TXRDY1
ENDIF
RET

PIN IF NORTHSTAR
IN DATA
ENDIF
IF APPLE
;MICRO IS APPLE
LDA DATA1
ENDIF
RET

```



```

      RET
EDOS1 EQU 0005H
EXX   EQU 0D9H
ATTN  EQU 52H
RXACK EQU 72H
DATA  EQU 04H
DATA1 EQU 0E09FH
TXRDY EQU 01H
TXRDY1 EQU 02H
RXRDY EQU 02H
RXRDY1 EQU 01H
STATUS EQU 05H
STATUS1 EQU 0E09EH
TXSYM  EQU 12H
TXACK  EQU 74H
RLDTA  EQU 0C3H
GOOD   EQU 67H
BAD    EQU 62H
DSKFUL EQU 64H
DMA    EQU 80H
ENDMA  EQU 01H
FCB    EQU 005CH
CTRLC  EQU 03H
CPN    EQU 0000H

;FLIP THE REGISTERS
;'R'
;'I'
;FCR NORTHSTAR
;FCR APPLE
;FCR NORTHSTAR
;FCR APPLE
;FCR NORTHSTAR
;FCR APPLE
;STATUS PORT FOR NORTHSTAR
;STATUS PORT FOR APPLE
;DC2 SYMBOL
;'t'
;OCBH MEANS DATA
;'g'
;'f'
;'d'
;ADDRESS OF DMA
;LAST LOCATION IN DMA
;ADDRESS OF FCB
;CCNTROL C MEANS GO TO CPM

```



```

GOON      EQU 47H      ;'G' MEANS CONTINUE
DONE      EQU 5AH      ;'Z' MEANS END OF SESSION
QUIT      EQU 14H      ;DC4 SYMBOL MEANS FILE COMPLETE
CONIN     EQU 01H      ;CHECK CONSOLE BUFFER FOR INPUT
CONOUT    EQU 02H      ;OUTPUT CURRENT A REG BYTE TO SCREEN
RIGHTS    DB 'MICROLAN  VERSION 2.0',13,10
DB 'COPYRIGHT (C) 1985  ROGER D. JASKOT and HAROLD W. HENRY$'
ENTER     DB 'ENTER NAME OF FILE TO BE SENT. IF THE FILE IS ON',13,10
DB 'A DISK IN THE OTHER DRIVE, ENTER IN THE FORMAT:',13,10,10
DB '      B:FILENAME.FILETYPE$'
WCHDSK    DB 'Write file to which disk drive?  Enter A for A drive,',13,10
DB 'B for B drive, or press return for default drive.$'
RXMODE    DB 'IN RECEIVE MODE.$'
FNFDMSG   DB 'FILE DOES NOT EXIST, RETURNING TO CPM.$'
TXING1    DB 'TRANSMITTING FILE.$'
HASFILE   DB 'RXING MICRO HAS FILE ALREADY, GOING TO CPM.$'
FULMSG    DB 'RXING MICRO DISK FULL.  RETURNING TO CPM.$'
WELCUM    DB 'WELCOME!',13,10,10
.
DB 'YOU ARE NOW ENTERING THE TRANSFER ZONE!$'
INSTRC    DB 'Enter an S for transmit mode, an R for receive mode,',13,10
DB 'or an X to exit.$'
FNFDMSG    DB 'FILE ALREADY EXISTS. RETURNING TO CPM.$'
EOFMSG    DB 'FILE TRANSMISSION COMPLETED.$'
NOMSG     DB 'NO FILE TRANSFER. RETURNING TO CPM.$'

```

```

ABRTMSG DB 'XMITTING MICRO ABORTED FILE TRANSFER.',13,10
        DB 'PLEASE ERASE FILENAME FROM YOUR DIRECTORY.$'
FULLMSG DB 'DISK FULL. FILE TRANSFER INCOMPLETE.$'
ERMSG   DB 'ENTER FILENAME AGAIN. END WITH <CR>$'
RXING1  DB 'CONNECTION MADE.$'
SWAP    DB EXX          ;SAVE THE REGISTERS USING EXX
        RET
CONBUF  DB 16          ;BUFFER FOR FILENAME
        DB 00
        DS 16

```

APPENDIX D

CP/M - 86

```
CSEG
ORG 0100H
MOV DX,OFFSET BAUDMSG      ;BAUDRATE HEADER
MOV CL,09H                 ;PFINT SAME
INT 0E0H
MOV CL,01H                 ;GET KEYBOARD INPUT
INT 0E0H
SUB AL,31H                 ;CCNVERT TO TABLE OFFSET
CMP AL,05H
JEE SETB1
JMP ERROR1

SETB1:  MOV BX,OFFSET TABL
ADD AL,AL
MOV AH,0
ADD BX,AX
MOV DX,[BX]
MOV BX,OFFSET BAUD
MOV [BX],DX
MOV DY,03FBH
MOV AL,83H
OUT DX,AL

;LINE CONTROL
;DIAB= 1
```

```

MOV DX,03F8H                ;BAUDATE DIVISOR
MOV EX,OFFSET BAUD
MOV AX,[BX]
OUT DX,AX
MOV DX,03FBH                ;CCNTROL
MOV AL,03H
OUT DX,AL                   ;RESET DLAB
INIT: CALL CRLF
MOV CL,09H                  ;PRINT STRING TO SCREEN
MOV DX,OFFSET RIGHTS        ;CCPYRIGHTS
CALL BDOS
CALL CRLF
MOV CL,09H                  ;PRINT STRING TO SCREEN
MOV DX,OFFSET WELCUM        ;WELCOME MSG
CALL BDOS
CALL CRLF
CALL CRLF
MOV CL,09H                  ;PRINT STRING TO SCREEN
MOV DX,OFFSET INSTRC        ;SEND, RECEIVE, OR QUIT?
CALL BDOS
CALL CRLF
CALL CRLF
HOLDING: MOV CL,06H          ;CHECK FOR CONSOLE INPUT
MOV DL,0FFH                 ;LOOKING FOR INPUT

```

```

CALL BDOS
AND AL,0DFH
CMP AL,53H
JNE G1 ! JMP MASTER
G1: CMP AL,52H
JNE G2 ! JMP SLAVE
G2: CMP AL,58H
JNE G3 ! JMP CPM
G3: JMP HOLDING
MASTER: MOV CL,09H
MOV DX,OFFSET ENTER
CALL BDOS
CALL CRLF
CALL CRLF
FILLUP: MOV BX,FCB
MCV [EX],BYTE PTR 00H
INC BX
MOV CL,0BH
FILLUP1: MOV [BX],BYTE PTR 20H ;FILL MEMORY ADDRESS WITH SPACES
INC BX
DEC CL
JNZ FILLUP1
MOV CL,13H
FILLUP2: MOV [BX],BYTE PTR 00H ;FILL REST OF ADDRESS WITH 0'S
;ENSURE LETTER IS A CAPITAL
;IS IT AN 'S'?
;IF SO, START FILE TRANSFER
;IS IT AN 'R'?
;IF SO, PREPARE TO RECEIVE FILE
;IS IT AN 'X'?
;IF SO, GO TO CPM
;REPEAT UNTIL INPUT FOUND
;PRINT STRING TO SCREEN
;ENTER FILENAME
;ADDRESS OF FCB
;11 SPACES
;MCVE PTR TO NEXT ADDRESS
;DECREMENT COUNTER
;REPEAT UNTIL DONE
;TOTAL OF 20 SPACES

```

```

INC BX
DEC CL
JNZ FILLUP2
HOLD1: MOV CL,0AH
MOV DX,OFFSET CCNBUF
CALL BDOS
MOV BX,OFFSET CCNBUF
MOV DX,FCB+1
INC BX
MOV CH,[BX]
MOV AL,[BX]
OR AL,AL
JNE G4 ! JMP ERROR
G4: INC BX
FLUP: MOV AL,[BX]
CMP AL,3AH
JNE G5 ! JMP DSKSEL
G5: CMP AL,2EH
JNE G6 ! JMP FIXIT
G6: CMP AL,40H
JNC G7 ! JMP DONTFIX
G7: AND AL,0DFH
DONTFIX: XCHG BX,DX ! MOV [BX],AL ! XCHG BX,DX ! STORE LETTER IN FCB
INC DX
;MCVE PTR TO NEXT ADDRESS
;DECREMENT COUNTER
;REPEAT UNTIL DONE
;READ CONSOLE BUFFER
;ADDRESS OF FIRST LETTER OF FILENAME
;ADDRESS OF CONSOLE BUFFER
;FCB ADDRESS
;STORE COUNT IN BX REGISTER
;MCVE COUNT TO ACCUMULATOR
;IS THERE AN INPUT?
;TRY AGAIN
;IS CHARACTER A ':'?
;IF SO, GO TO DISK SELECT
;IS IT A '.'?
;IF SO, SKIP TO FILETYPE
;CHECK FOR LETTER
;SKIP NEXT STEP IF NOT LETTER
;ENSURE LETTER IS A CAPITAL
;AL ! XCHG BX,DX ! STORE LETTER IN FCB

```



```

INC BX
DEC CH
JNZ FLUP
LISN1: MOV AL,ATTN
CALL POUT1
MOV CL,03H
LISN: CALL PIN
CMP AL,RXACK
JNE G8 ! JMP XMIT
G8: DEC CL
JNZ LISN
JMP LISN1
XMIT: CALL CRLF
CALL CRLF
MOV CL,09H
MOV DX,OFFSET RXING1
CALL BDOS
CALL CRLF
XMIT1: MOV AL,TXSYM
CALL POUT1
MOV CL,08AH
LITTLT: CALL PIN
CMP AL,TXACK
JNE G9 ! JMP TXFCE
;REPEAT UNTIL END
;LETTER 'R'
;LISTEN 3 TIMES
;IS IT AN 'r'?
;IF SO, THEN XMIT
;OTHERWISE DECREMENT CTR
;LISTEN UNTIL CTR IS ZERO
;THEN TRY AGAIN
;PRINT STRING TO SCREEN
;'r' WAS RECEIVED
;DC2 SYMBOL FOR SYNC AT START
;OF 128 BYTE BLOCK
;LISTEN 138 TIMES
;WAS 't' RECEIVED?
;IF SO, XMIT FILE CTRL BLK

```

```

G9: DEC CL
    JNZ IITTLT
    JMP XMIT1
TXFCB: CALL CRLF
    CALL OPENIT
TXFCB1: MOV CH,00H
    MOV BX,FCB
FCBLUP: MOV AL,CH
    INC BX
    XOR AL,[BX]
    MOV CH,AL
    MOV AL,[BX]
    CALL POUT1
    CMP AL,0H
    JNE G10 : JMP FCBCX
G10: JMP FCBLUP
FCBCX: MOV CL,20H
FCBCX1: CALL PIN
    CMP AL,RXACK
    JNZ FCBCX1
    MOV AL,CH
    CALL POUT
    PUSH CX
    MOV AL,0H
;CITHERWISE KEEP LISTENING
;UNTIL CTR IS ZERO,
;THEN SEND DC2 SYNC AGAIN
;SEND FILENAME TO RXING MICRO
;SEE IF FILE EXISTS. IF SO, OPEN IT
;INITIALIZE CHECKSUM REGISTER
;SET PTR TO 1ST LETTER IN FILENAME
;PERFORM CHECKSUM OPERATION
;MOVE PTR TO NEXT BYTE
;BY XORING CURRENT BYTE
;WITH B REGISTER
;PUT CURRENT BYTE IN ACCUM
;SEND CURRENT BYTE
;CHECK FOR END OF FILENAME
;IF END, GO TO CHECKSUM LOOP
;IF NOT, REPEAT FCB LOOP
;LCOP 32 TIMES
;FCR SYNC WITH SLAVE
;IS IT AN 'r'
;IF NOT, LISTEN AGAIN. IF SO,
;PUT CHECKSUM IN ACCUM
;SEND CHECKSUM
;SAVE CHECKSUM
;CLEAR ACCUM

```

```

MOV      CH,80H
FCB TMCT: CALL PIN
      CMP  AL,BAD
      JNE G11 ! JMP RSNDFCB
G11:  CMP  AL,GOOD
      JNE G12 ! JMP WAITFIL
G12:  DEC  CH
      JNZ  FCB TMOT
      PCP  CX
      DEC  CL
      JNZ  FCBCK1
      JMP  TXFCB
      WAITFIL: POP CX
      MOV  CX,07FFH
      WAIT1: CALL STATIN1
      JZ   WAIT1
      DEC  CX
      JNE G13 ! JMP GOCPM
G13:  CALL PIN
      CMP  AL,QUIT
      JNE G14 ! JMP GOCPM1
G14:  CMP  AL,GOON
      JNZ  WAITFIL
;LISTEN 100 TIMES
;READ MAIL
;DID IT CHECK BAD?
;IF SO, SEND FCB AGAIN
;DID IT CHECK GOOD?
;IF SO, GO TO NEXT ROUTINE
;IF NOT, DECREMENT CTR, AND
;IF NOT 0, LISTEN AGAIN
;CLEAR STACK
;IF SO, DECREMENT CL
;AND REPEAT UNTIL CL=0
;IF 0, ASSUME PROBLEM AND SEND AGAIN
;CLEAR STACK
;CCUNT LOOP APPX 2K
;ANY 'MAIL'?
;IF NOT, CHECK AGAIN
;IF SO, DECREMENT CTR
;AND, IF 0, QUIT
;OTHERWISE READ 'MAIL'
;DCES RXING MICRO ALREADY HAVE FILE?
;IF SO, GO TO CPM
;IS IT THE GO ON SIGNAL 'G'
;IF NOT, LISTEN AGAIN. ALLOW RXING
;MICRO TO CATCH UP

```

```

CALL CRLF
TXDATA:
MOV CL,09H
MOV DX,OFFSET TXING1
CALL BDOS
CALL CRLF
RDSEQ: CALL READSEQ
SEND: CALL CHECK
MOV AL, TXSYM
CALL POUT1
MOV CL,0FH
LITLET2: CALL PIN
CMP AL, TXACK
JNE G15 : JMP SLUP2
G15: CMP AL, DSKFUL
JNE G16 : JMP FULDISK
G16: DEC CL
JNZ LITLET2
JMP SEND
SLUP2: MOV AL, RLDTA
CALL POUT1
MOV CX, 07FFH
SLUP3: CALL PIN
CMP AL, RLDTA
;SEND THE FILE
;PRINT STRING TO SCREEN
;SAYS FILE BEING SENT
;READ FIRST 128 BYTE BLOCK
;AND SEND TO RXING MICRO
;DC2 SYMBOL FOR SYNC AT START OF DATA
;LISTEN 15 TIMES
;IS IT A 't'?
;IF SO, READY TO SEND DATA
;IS RXING MICRO'S DISK FULL?
;IF SO, QUIT
;IF NOT, DECREMENT CTR
;LISTEN AGAIN, UNLESS CTR IS 0,
;THEN TRY TO SYNC AGAIN
;OCBH MEANS TIME FOR DATA
;WAIT LOOP APPX 2K
;IS IT AN ECHO?

```

```

JZ SLUP1
DEC CX
JNZ SLUP3
JMP SLUP2

SLUP1: MOV BX,DMA
MOV CH,00H
SLOOP: MOV AL,CH
XOR AL,[BX]
MOV CH,AL
MOV AL,[BX]
CALL POUT
INC BX
MOV AL,BH
CMP AL,ENDMA
JNZ SLOOP

CRC: MOV AL,CH
CALL POUT

CRCTMCT: MOV CH,01AH
CRCT1: CALL STATIN1
JZ CRCT1
CALL PIN
CMP AL,BAD
JNE G17 ! JMP RESEND
G17: CMP AL,GOOD

;IF SO, SEND DATA
;IF NOT, DECREMENT COUNTER
;REPEAT UNTIL ZERO

;PCINTER TO 1ST INFO BYTE
;INITIALIZE CHECKSUM LOCATION
;PERFORM CHECKSUM
;XCR DATA WITH CH REGISTER

;PUT BYTE IN ACCUMULATOR
;DATA IS TRANSFERRED
;MCVE PTR TO NEXT BYTE
;MCVE BH REG TO ACCUM
;END DMA, CHECK FOR LAST EYTE
;IF NOT, SEND NEXT BYTE. OTHERWISE
;PUT CHECKSUM IN ACCUMULATOR
;AND SEND TO RXING MICRC
;LISTEN 26 TIMES
;CHECK INPUT BUFFER
;IF NOTHING, TRY AGAIN
;READ MAIL
;IS CHECK BAD?
;IF SO, SEND BLOCK AGAIN
;IS CHECK GOOD?

```

```

JNE G18 ! JMP RDSQRPT
G18: DEC CH
JNZ CRCT1
JMP SEND

DSKSEL: MOV BX,OFFSET CONBUF+2 ;ADDRESS OF DISK SEL ENTRY
MOV AL,[BX]
AND AL,ODFH
CMP AL,'A'
JZ ADISK
CMP AL,'B'
JZ PDISK
CMP AL,'C'
JZ CDISK
JMP DSKSEL1

ADISK: MOV DI,FCB
MOV [DI],BYTE PTR 01H
JMP DSKSEL1

BDISK: MOV DI,FCB
MOV [DI],BYTE PTR 02H
JMP DSKSEL1

CDISK: MOV DI,FCB
MOV [DI],BYTE PTR 03H
JMP DSKSEL1

;IF SO, READ NEXT BLOCK
;DECREMENT COUNTER
;IF NOT TIMED OUT,LISTEN AGAIN
;IF TIMED OUT, ASSUME PROBLEM.
;SEND BLOCK AGAIN
;PUT DISK SEL IN ACCUM
;ENSURE LETTER IS CAPITAL
;IS LETTER AN 'A'?
;IF SO, SET FOR A DRIVE.
;IS LETTER A 'B'?
;IF SO, SET FOR B DRIVE.
;IS LETTER A 'C'?
;IF SO, SET FOR C DRIVE.
;IF NEITHER, RETURN TO FILENAME LOOP.
;SET PTR TO DRIVE BYTE.
;SET FCB FOR A DRIVE.
;RETURN TO FILENAME LOOP.
;SET PTR TO DRIVE BYTE.
;SET FCB FOR B DRIVE.
;RETURN TO FILENAME LOOP.
;SET PTR TO DRIVE BYTE.
;SET FCB FOR C DRIVE.
;MOVE BUFFER POINTER TO FILENAME.

```



```

INC BX
MOV DX,FCB+1
JMP FLUP
RESEND: MOV CL,CONOUT
MOV DL,BAD
CALL BDOS
JMP SEND
RSNDFCB: POP CX
MOV AL,CH
ADD AL,3
CALL POUT1
RSNDF: CALL PIN
CMP AL,RXACK
JNZ RSNDF
RSNDFC1: CALL PIN
CMP AL,TXACK
JNZ RSNDFC1
JMP TXFCB1
FIXIT: MOV DX,FCB+9
INC BX
JMF FLUP
ERROR: MOV CL,09H
MOV DX,OFFSET FMSG
CALL BDOS
;FCB FILENAME ADDRESS.
;RETURN TO FILENAME LOOP.
;PRINT TO SCREEN
;A 'b' IF CHECKSUM WAS BAD
;AND SEND BLOCK AGAIN
;RECALL CHECKSUM
;PUT CHECKSUM IN ACCUM
;ADD 3 TO OFFSET
;SEND BYTE
;IS IT AN 'r'
;IF NOT LISTEN AGAIN
;READ MAILBOX
;SYNC WITH RXING MICRO
;REPEAT UNTIL TXACK RECEIVED
;IF SO, RESEND FCB
;MCVE POINTER TO FILETYPE AREA
;MCVE PTR TO FIRST LETTER OF FILTYPE
;PRINT STRING TO SCREEN
;ERROR MESSAGE

```

```

CALL CRLF
JMP HOLD1
POUT1: LAHF ! PUSH AX

        MOV CL,06H
        MOV DL,0FFH
        CALL BDOS
        CMP AL,CTRLC
        JNE G22 ! JMP STOPS
G22: CALL CHECK
        POP AX ! SAHF
        MOV DX,DATA
        OUT DX,AL
        RET

OPENIT: MOV CL,0FH
        MOV DX,FCB
        CALL BDOS
        CMP AL,0FFH
        JNE G23 ! JMP FNFOUND
G23: RET
CLOSIT: MOV CL,10H
        MOV DX,FCB
        CALL BDOS
CLOSIT1: MOV AL,DONE

        ;LOOK FOR INPUT AGAIN
        ;SEND THE DATA
        ;FIRST, SAVE THE CURRENT BYTE
        ;CHECK FOR CONSOLE INPUT
        ;LOOKING FOR INPUT

        ;IS THERE A CONTROL C?

        ;PERFORM CHECK
        ;AND RECALL BYTE

        ;OPEN FILE CODE
        ;FILE CTRL BLOCK ADDRESS IN DX REG PR

        ;FF = FILE NOT FOUND
        ;IF FILE NOT FOUND
        ;CTHERWISE, RET TO TX DATA
        ;CLOSE FILE CODE
        ;FILE CTRL BLOCK ADDRESS IN DX REG PR

        ;END OF SESSION MSG 'Z'

```

```

CALL POUT1
MOV AL,0H
CALL PIN
CMP AL,DONE
JNZ CLOST1
JMP GOCPM
READSEQ: PUSH CX
        PUSH DX
        MOV CL,14H
        MOV DX,FCB
        CALL BDOS
        POP DX
        POP CX
        CMP AL,0
        JZ G24 ! JMP EOF1LE
G24: RET
RDSQRPT: MOV CL,CONCUT
        MOV DL,02AH
        CALL BDOS
        JMP RDSEQ
FNFOUND: MOV AL,QUIT
        CALL POUT1
        MOV CL,09H
        MCV DX,OFFSET FNFDMG
;SEND TO RXING MICRO
;CLEAR ACCUM
;CHECK REPLY
;DCES RXING MICRO AGREE?
;IF NOT, REPEAT
;IF SO, GO TO CPM
;READ SEQUENTIAL CODE
;FILE CTRL BLOCK ADDRESS IN DX REG PR
;0 MEANS SUCCESSFUL READ
;IF NOT 0, ASSUME FINISHED WITH FILE
;PRINT TO SCREEN
;'' SO USER KNOWS BLK WAS SENT
;TC READ NEXT 128 BYTE BLK
;TELL RXING MICRO NO FILE FOUND
;PRINT STRING TO SCREEN
;FILE NOT FOUND MSG

```

```

CALL BDOS
CALL CRLF
JMP GOCPM

EOFIL: MOV AL, TXSYM
CALL POUT1
MOV CL, 0FH
LITLET3: CALL STATIN1
JZ LITLET3
CALL PIN
CMP AL, TXACK
JNE G25 ! JMP EOFIL1

G25: DEC CL
JNZ LITLET3
JMP EOFIL

EOFIL1: MOV AL, QUIT
CALL POUT1
CALL PIN
CMP AL, QUIT
JNZ EOFIL1
CALL CRLF
MOV CL, 09H
MOV DX, OFFSET ECFMSG
CALL BDOS
CALL CRLF

;AND GO TO CPM
;DC2 SYMBOL FOR SYNC WITH RXING MICRO

;LISTEN 15 TIMES
;CHECK FOR MAIL
;IF NONE, CHECK AGAIN
;READ MAIL
;IS IT A 't'?
;IF SO, CONTINUE
;IF NOT, DECREMENT COUNTER
;AND LISTEN AGAIN, UNLESS COUNTER IS
;0. THEN TRY AGAIN
;DC4 SYMBOL. TELLS RXING MICRO THAT
;THE FILE IS DONE
;LISTEN FOR REPLY
;DCES RXING MICRO ACKNOWLEDGE?
;IF NOT, TRY AGAIN

;PRINT STRING TO SCREEN
;IF SO, TELL USER FILE IS DONE

```

```

JMP CLOST
STOPS: MOV AL,CTRLC
MOV DX,DATA
OUT DX,AL
MOV AL,0H
CALL PIN
CMP AL,CTRLC
JNZ STOPS
JMP GOCPM

FULDISK: MOV AL,DONE
CALL POUT1
MOV CL,09H
MOV DX,OFFSET FULMSG
CALL BDOS
CALL CRLF
JMP GOCPM

GOCPM: MOV AL,0H
MOV DX,DATA
OUT DX,AL
CALL CRLF
CALL PIN
JMP CPM

GOCPM1: MOV CL,09H
MOV DX,OFFSET HASFILE
;AND CLOSE THE FILE
;SEND CTRLC TO RXING MICRO
;CLEAR ACCUM
;FROM RXING MICRO
;ACK FROM RXING MICRO
;REPEAT UNTIL ACK
;LETTER 'Z' TO ACKNOWLEDGE
;SEND BYTE
;PRINT STRING TO SCREEN
;SAYS RXER'S DISK FULL
;RESET THE ACCUMULATOR AND
;CLEAR OUTPUT BUFFER
;AND GO TO CPM
;PRINT STRING TO SCREEN
;RXING MICRO HAS FILE ALREADY

```

```

CALL BDOS
JMP GOCPM
SLAVE: MOV CL,09H
MOV DX,OFFSET WCHDSK
CALL BDOS
CALL CRLF
CALL CRLF
DRVSEL: MOV CL,06H
MOV DL,0FFH
CALL BDOS
CMP AL,0DH
JNE G26 ! JMP CONT
G26: AND AL,0DFH
CMP AL,'A'
JZ G27 ! JMP DISK3
G27: MOV BX,FCB
MOV [BX],BYTE PTR 01H
JMF CONT
DISKB: CMP AL,'B'
JNZ DISKC
MOV BX,FCB
MOV [EX],BYTE PTR 02H
JMP CONT
DISKC: CMP AL,'C'

;PRINT STRING TO SCREEN
;SELECT DISK DRIVE

;CHECK FOR CONSOLE INPUT
;LOOKING FOR INPUT

;IS IT A <CR>?
;IF SO, ENTER RECEIVE MODE
;ENSURE LETTER IS A CAPITAL
;IS IT AN 'A'?
;SKIP TO B IF NOT 'A'
;ADDRESS OF DISK DRIVE BYTE
;SET BYTE TO A DISK DRIVE
;THEN CONTINUE
;IS IT A 'B'?
;IF NOT, SKIP TO C
;ADDRESS OF DISK DRIVE BYTE
;SET BYTE TO B DISK DRIVE
;THEN CONTINUE
;IS IT A 'C'?

```



```

JNZ DRVSEL
MOV BX,FCB
MOV [BX],BYTE PTR 03H
CONT: MOV CL,09H
MOV DX,OFFSET EXMODE
CALL BDOS
CALL CRLF
SLAVE1: MOV AL,00H
CALL PIN
CMP AL,ATTN
JNZ SLAVE1

CALL CRLF
MOV CL,09H
MOV DX,OFFSET RXING1
CALL BDOS
CALL CRLF
MOV AL,RXACK
CALL POUT
LISTEN: CALL PIN1
CMP AL,TXSYM
JNZ LISTEN
CALL CRLF

; IF NOT, LISTEN AGAIN
; ADDRESS OF DISK DRIVE BYTE
; SET BYTE TO C DISK DRIVE
; PRINT STRING TO SCREEN
; IN RECEIVE MODE

; RESET ACCUMULATOR
; LISTENING FOR AN 'R'
; 'F'
; IF 'R' RX'D, CONTINUE. IF NOT
; LISTEN AGAIN

; PRINT STRING TO SCREEN
; CONNECTION MADE

; 'r'
; SEND AN 'r' TO XMITTING MICRO
; LISTENING FOR A 'DC2'
; 'DC2'
; IF 'DC2' RX'D, CONTINUE. IF NOT,
; LISTEN AGAIN
CALL CRLF

```

```

RYFCB: MOV  BX,FCB+1      ;ADDRESS OF FCB MEM LOC INTO BX REG PR
      MOV  CL,1EH         ;COUNTER FOR FCB'S 31 SPACES
RSTFCB: MOV  [BX],BYTE PTR 00H ;FILL FCB WITH 0'S
      INC  BX             ;MCVE PTR TO NEXT MEMORY ADDRESS IN FCB
      DEC  CL             ;DECREMENT COUNTER
      MOV  AL,CL
      CMP  AL,0
      JNZ  RSTFCB

      MOV  CH,00H
      MOV  BX,FCB+1
      MOV  AL,TXACK
      CALL POUT
      CALL PIN1
RST1:  CALL STATIN1
      JZ   RST1

RST2:  CALL PIN1
      CMP  AL,QUIT
      JNE G28 ! JMP NOFILE
G28:   CMP  AL,0H
      JNE G29 ! JMP FCECRC
G29:   CMP  AL,TXSYN
      JZ   RST2
      MOV  [BX],AL

      ;IF COUNTER = 0, CONT. IF NOT,
      ;PUT ANOTHER 0 IN FCB
      ;INITIALIZE CHECKSUM
      ;ICAD 2ND ADDRESS OF FCB IN
      ;'t'
      ;SEND 't' TO XMITING MICRC FOR SYNC
      ;CLEAR THE ACCUMULATOR
      ;CHECKING FOR INPUT

      ;FILE NAME DATA
      ;IS DATA A 'QUIT'?
      ;FILE DID NOT EXIST
      ;CHECK IF FILENAME COMPLETELY SENT
      ;IF FILENAME RX'D, GO TO CHECKSUM
      ;CHECK IF DATA IS VALID
      ;IF DATA IS NOT FILENAME,
      ;PUT FILENAME IN FCB

```

```

CALL OUTPUT
MOV AL,CH
XOR AL,[BX]
MOV CH,AL
INC BX
JMP RST1

FCBCRC: MOV AL,RXACK
CALL POUT
FCBCRC1: CALL STATIN1
JZ FCBCRC1
CALL CRLF
CALL PIN
CMP AL,CH
JNE G30 ! JMP STRIFIL
G30: ADD AL,3
MOV CL,AL
MOV AL,BAD
CALL POUT
CLEAR: CALL PIN
CMP AL,CL
JNZ CLEAR
MOV AL,RXACK
CALL POUT
JMP RXFCB

;PRINT FILENAME TO SCREEN

;CALCULATE CHECKSUM

;MOVE PTR TO NEXT FCB ADDRESS

;'I'
;SYNC DATA WITH XMITING MICRO
;CHECKING FOR INPUT

;CHECKSUM DATA
;CCMPARE CHECKSUM
;CHECKSUM MATCHED
;AID 3 TO THE CHECKSUM
;STORE IN REGISTER
;CHECKSUM DID NOT MATCH
;TELL XMITING MICRO

;XMITING MICRO STOPPED SENDING CHKSUM?
;IF NOT,LISTEN AGAIN
;SYNC WITH XMITING MICRO

;TRY AGAIN

```

```

STRTRFIL: MOV AL,GOOD
CALL POUT
CALL OPNFILE
CALL MAKEFIL
RXD1: MOV CH,00H
MOV BX,DMA
MOV CL,81H
RXDS: CALL STATIN1
JZ RXDS
RXD2: CALL PIN1
CMP AL,TXSYM
JNZ RXD2
MOV AL,TXACK
CALL POUT
RXDS1: CALL STATIN1
JZ RXDS1
RXYT: CALL PIN1
CMP AL,RDLTA
JZ RXYET1
CMP AL,QUIT
JNZ RXYET4
JMP CLSFILE
RXYET4: JMP RXDS1

;READY TO CHECK IF FILE ALREADY PRESENT
;CHECK IF FILE EXISTS
;CREATE NEW FILE
;INITIALIZE CHECKSUM
;LCAD ADDRESS OF DMA MEM LOC TO
;EX REGISTER PAIR
;INITIALIZE COUNTER WITH SIZE OF DMA

;SYNC WITH XMITING MICRO
;CCMPARE WITH 'DC4'
;
;t'
;IN SYNC WITH XMITING MICRO
;CHECKING FOR INPUT

;IS IT OCBH?
;IF SO, GO TO RECEIVF DATA
;IS IT 'DC4' FOR QUIT?
;IF NOT, LISTEN AGAIN; OTHERWISE,
;IF SO, CLOSE FILE
;LISTEN AGAIN

```

RXVET1: MOV AL,RDLTA CALL POUT MOV AL,00H RXVET2: CALL STATIN1 JZ RXVET2 RXVET3: CALL PIN CMP AL,RDLTA JZ RXVET3 RXD3: DEC CL JNE G32 ! JMP RXCRC G32: MOV [BX],AL MOV AL,CH XOR AL,[BX] MOV CH,AL INC BX RXD4: CALL STATIN1 JZ RXD4 CALL PIN JMP RXD3 EXCRC: MOV AL,AL CMP AL,CH JNE G33 ! JMP WRITFIL G33: MOV AL,BAD CALL POUT	;ACK REAL DATA COMING ;CLEAR ACCUM ;CHECKING FOR INPUT ;READ DATA ;IS IT STILL RDLTA? ;DECREMENT COUNTER ;CHECKSUM RX'D ;PUT THE DATA IN MEMORY ;CALCULATE CHECKSUM ;MCVE PTR TO NEXT DMA ADDRESS ;CHECK FOR INPUT ;ICOP UNTIL INPUT ;ENSURE CH IS COMPARED TO A ;COMPARE WITH CHECKSUM ;128 BYTE BLOCK SENT ;CHECKSUM DID NOT MATCH ;NOTIFY XMITING MICRO
--	--

```

JMP  RXD2
POUT: LAHF ! PUSH AX
      CALL CHECK
      PCP AX ! SAHF
      MOV DX,DATA
      OUT DX,AL
      RET
PIN1: MOV DX,DATA
      IN AL,DX
      CMP AL,CTRLC
      JNE G34 ! JMP ABORT
G34:  RET
WRITEFIL: MOV  AL,GOOD
      CALL POUT
      CALL WRITSEQ
      MOV  CL,CONOUT
      MOV  DL,02AH
      CALL BDOS
      JMP  RXD1
OPNFILE: MOV  CL,0FH
      MOV  DX,FCB
      CALL BDOS
      CME  AL,0FFH
      JZ  G35 ! JMP FILFND
;SEND 128 BYTE BLOCK AGAIN
;SAVE THE DATA
;RETURN THE DATA
;SEND DATA
;DID XMITING MICRO ABORT?
;IF SO, ABORT
;XMIT THAT THE CHECKSUM IS CORRECT
;START WRITING FILE TO DISK
;PRINT TO SCREEN
;'*' TO PRINT TO SCREEN
;OPEN FILE CODE
;FCB ADDRESS IN DX RGSTF PAIR
;FF = FILE NOT FOUND
;FILE EXISTS

```



```

G35: RET
CLSFILE: MOV AL,QUIT
CALL POUT
MOV CL,10H
MOV DX,FCB
CALL BDOS
CALL CRLF
MOV DX,OFFSET ECFMSG
MOV CL,09H
CALL BDOS
CALL CRLF
MOV AL,0H
CLSFILE1: CALL PIN1
CMP AL,DONE
JNZ CLSFIL1
MOV AL,DONE
CALL POUT
JMF CPM
MAKEFIL: MOV CL,16H
MOV DX,FCB
CALL BDOS
MOV AL,GOON
CALL POUT
RET
; 'IC4'
; AGREE END OF FILE
; CLOSE FILE CODE
; FCB ADDRESS IN DX RGSTR PAIR
; FILE TRANSMISSION COMPLETED
; PRINT STRING TO SCREEN
; CLEAR THE ACCUMULATOR
; LOOKING FOR END OF SESSION MSG
; 'Z' = END OF SESSION
; END OF SESSION MESSAGE
; CCNFRM RECEPTION OF E-O-SESSION MSG
; MAKE NEW FILE CODE
; FCB ADDRESS IN DX RGSTR PAIR
; CCNTINUE MESSAGE
; RETURN TO RX FIRST 128 BYTE BLOCK

```

```

WKITSEQ: PUSH CX
        PUSH DX
        MOV  CL,15H
        MOV  DX,FCB
        CALL BDOS
        POP  DX
        PCP  CX
        OR   AL,AL
        JZ  G36 ! JMP FULLDSK
G36: RET
FILEND: MOV  AL,QUIT
        CALL POUT
        MOV  CL,09H
        MOV  DX,OFFSET FNDMSG
        CALL BDOS
        CALL CRLF
        JMF  CPM
NOFILE: MOV  CL,09H
        MOV  DX,OFFSET NCMSG
        CALL BDOS
        CALL CRLF
        JMF  CPM
ABORT: CALL CRLF
        MOV  AL,CTRLC
        SEND XMITTING MICRO ABORT ACK
;WRITE THE FILE TO THE DISK
;FCB IN DX RGSTR PAIR
;CHECK IF DISK IS FULL
;IF SO, JUMP TO FULLDSK
;TELL XMITTING MICRO, FILE FOUND
;PRINT STRING TO SCREEN
;FILE ALREADY EXISTS. GO TO CPM
;PRINT STRING TO SCREEN
;NC FILE TRANSFER

```

```

CALL POUT
MOV CL,09H
MOV DX,OFFSET AERTMSG
CALL BDOS
CALL CRLF
JMP GOCPM
FULLDSK: MOV AL,DSKFUL
CALL POUT
CALL PIN
CMP AL,DONE
JNZ FULLDSK
MOV CL,09H
MOV DX,OFFSET FULLMSG
CALL BDOS
CALL CRLF
JMP CPM
CHECK: CALL STATIN2
JZ CHECK
RET
STATIN1: MOV DX,STATUS
IN AL,DX
AND AL,RXRDY
RET
STATIN2: MOV DX,STATUS

```

```

;PRINT STRING TO SCREEN

```

```

;XMITTING MICRO ABORTED

```

```

;GC TO CPM

```

```

;'L'

```

```

;TELL XMITTING MICRO DISK FULL

```

```

;AWAITING CONFIRMATION

```

```

;PRINT TO SCREEN

```

```

;FILE TRANSFER INCOMPLETE, DISK FULL

```

```

;CHECK STATUS BYTE

```

```

;CCNTINUE UNTIL TXRDY IS SET

```

```

IN AL,DX
AND AL,TRDY
RET
PIN: MOV DX,DATA
IN AL,DX
RET
CRLF: PUSH AX
MOV AL,ODH
CALL OUTPUT
MOV AL,0AH
CALL OUTPUT
POP AX
RET
OUTPUT: PUSH ES
PUSH BX
PUSH DX
PUSH CX
LAHF ! PUSH AX
MOV CL,CONOUT
MOV DI,AL
CALL BDOS
POP AX ! SAHF
POP CX
POP DX
;CARRIAGE RETURN
;LINE FEED
;SAVE THE ES
;BX,
;DX,
;AND CX REGISTERS
;PRINT TO SCREEN
;PUT THE ACCUMULATOR IN 'DI' RGSTR
;RETURN THE CX,
;DX,

```

```

POP BX
POP ES
RET
;BX
;AND ES REGISTERS

```

```

BDOS: PUSH ES
      PUSH BX
      PUSH DX
      PUSH CX
      INT 0E0H
      POP CX
      POP DX
      PCP BX
      POP ES
      RET
;SAVE THE ES,
;BX,
;DX,
;AND CX REGISTERS
;EXECUTE
;RETURN THE CX,
;DX,
;PX
;AND ES REGISTERS

```

```

CPM: MOV DL,00H
      MOV CL,00H
      JMP BDOS
;RETURN TO OPERATING SYSTEM

```

```

ERROR1: MOV DX,OFFSET ERR1
        MOV CL,09H
        ;BAUDRATE ERROR
        ;PRINT STRING TO SCREEN

```

```

      INT 0E0H
      MOV CL,0
      MOV DL,0
      INT 0E0H

```

```

ATTN EQU 52H
RXACK EQU 72H
; 'I'
; 'I'

```

```

DATA      EQU 03F8H
TXRDY     EQU 20H
RXRDY     EQU 01H
STATUS    EQU 03FDH
TXSYM     EQU 12H
TXACK     EQU 74H
RLDTA     EQU 0CBH
GOOD      EQU 67H
BAD       EQU 62H
DSKFUL    EQU 64H
DMA       EQU 80H
ENDMA     EQU 01H
FCB       EQU 005CH
CTRLC     EQU 03H
GOON      EQU 47H
DONE      EQU 5AH
QUIT      EQU 14H
CONIN     EQU 01H
CONOUT    EQU 02H
RIGHTS    DB 'MICROLAN  VERSION 2.1',13,10
ENTER     DB 'COPYRIGHT (C) 1985  ROGER D. JASKOT AND HAROLD W. HENRY$'
          DB 'ENTER NAME OF FILE TO BE SENT. IF THE FILE IS ON',13,10
          DB 'A DISK IN ANOTHER DRIVE, ENTER IN THE FORMAT:',13,10,10
          DE ' B:FILENAME.FILETYPE$'

```



```

WCHDSK DB 'WRITE FILE TO WHICH DISK DRIVE? ENTER AN A FOR A DRIVE,'13,10
        DB 'A B FOR E DRIVE,... OR PRESS RETURN FOR DEFAULT DRIVE.$'
RXMODE DB 'IN RECEIVE MODE.$'
FNFDMG DB 'FILE DOES NOT EXIST, RETURNING TO CPM.$'
TXING1 DB 'TRANSMITTING FILE.$'
HASFILE DB 'RXING MICRO HAS FILE ALREADY, GOING TO CPM.$'
FULMSG DB 'RXING MICRO DISK FULL. RETURNING TO CPM.$'
WELCUM DB 'WELCOME!',13,10,10
INSTRC DB 'YOU ARE NOW ENTERING THE TRANSFER ZONE!$'
        DB 'ENTER AN S FOR TRANSMIT MODE, AN R FOR RECEIVE MODE,'1,13,10
        DB 'OR AN X TO EXIT.$'
FNDMSG DB 'FILE ALREADY EXISTS. RETURNING TO CPM.$'
FOFMSG DB 'FILE TRANSMISSION COMPLETED.$'
NOMSG DB 'NO FILE TRANSFER. RETURNING TO CPM.$'
ABRTMSG DB 'XMITING MICRO ABORTED FILE TRANSFER.',13,10
        DB 'PLEASE ERASE FILENAME FROM YOUR DIRECTORY.$'
FULLMSG DB 'DISK FULL. FILE TRANSFER INCOMPLETE.$'
ERMMSG DB 'ENTER FILENAME AGAIN. END WITH <CR>$'
RXING1 DB 'CONNECTION MADE.$'
CONBUF DB 16 ;BUFFER FOR FILENAME
        DB 00
        RS 16
BAUDMSG DB 'SELECT PAUD RATE',CDH,0AH
        DB '1 = 300 BAUD',ODH,0AH

```

```

DB '2 = 600  BAUD',0DH,0AH
DB '3 =1200  BAUD',0DH,0AH
DB '4 =2400  BAUD',0DH,0AH
DB '5 =4800  BAUD',0DH,0AH
DB '6 =9600  BAUD$'

```

```

ERR1 DE 'BAUDRATE OUT OF RANGE$'

```

```

BAUD RW 1

```

```

TABL DB 80H,01,0C0H,00,60H,00,3AH,00,18H,00,0CH,00

```

NOTE: This version of MICROLAN was obtained by direct conversion through the TRANS86 conversion program. We did not attempt to streamline the procedures or efficiency of operation for this IBM compatible version.

APPENDIX E

MS.DOS

```
STACK      SEGMENT PARA STACK 'STACK'
DB 512 DUP(0)
STACK      ENDS
DATA       SEGMENT PARA PUBLIC 'DATA'
DTA        DB 80H DUP(0)
FCB        DB 36 DUP(0)
ATTN       EQU 52H      ;'F'
RXACK      EQU 72H      ;'I'
DATA 1     EQU 03F8H
TXRDY      EQU 20H
RXRDY      EQU 01H
STATUS     EQU 03FDH
TXSYM      EQU 12H      ;DC2 SYMBOL
TXACK      EQU 74H      ;'t'
RLDTA      EQU 0CBH
GOOD       EQU 67H      ;'c'
EAD         EQU 62H      ;'E'
DSKFUL     EQU 64H      ;'d'
DMA         EQU DTA      ;ADDRESS OF DMA
ENDMA      EQU 01H      ;LAST LOCATION IN DMA
CTRLC      EQU 03H      ;CONTROL C MEANS GO TO CPM
```

```

GOON      EQU 47H      ;'C' MEANS CONTINUE
DONE      EQU 5AH      ;'Z' MEANS END OF SESSION
QUIT      EQU 14H      ;DC4 SYMBOL MEANS FILE COMPLETE
CONIN     EQU 01H      ;CHECK CONSOLE BUFFER FOR INPUT
CONOUT    EQU 02H      ;OUTPUT CURRENT A REG BYTE TO SCREEN
RIGHTS    DB 'MICROLAN  VERSION 2.2',13,10
DB 'COPYRIGHT (C) 1985  ROGER D. JASKOT AND HAROLD W. HENRY$'
ENTER     DB 'ENTER NAME OF FILE TO BE SENT. IF THE FILE IS ON',13,10
DB 'A DISK IN THE OTHER DRIVE, ENTER IN THE FORMAT:',13,10,10
DB '      B:FILENAME.FILETYPE$'
WCHDSK    DE 'WRITE FILE TO WHICH DISK DRIVE?  ENTER AN A FOR A DRIVE,',13,10
DB 'B FOR B DRIVE..., OR PRESS RETURN FOR DEFAULT DRIVE.$'
RXMODE     DB 'IN RECEIVE MODE.$'
PWFDMMSG  DB 'FILE DOES NOT EXIST, RETURNING TO CPM.$'
TXING1    DB 'TRANSMITTING FILE.$'
HASFILE    DB 'RXING MICRO HAS FILE ALREADY, GOING TO CPM.$'
FULMSG    DB 'RXING MICRO DISK FULL.  RETURNING TO CPM.$'
WFICUM     DB 'WELCOME!',13,10,10
DB 'YOU ARE NOW ENTERING THE TRANSFER ZONE$'
INSTRC     DB 'ENTER AN S FOR TRANSMIT MODE, AN R FOR RECEIVE MODE,',13,10
DB 'OR AN X TO EXIT.$'
ENDMSG     DB 'FILE ALREADY EXISTS. RETURNING TO CPM.$'
EOFMSG     DB 'FILE TRANSMISSION COMPLETED.$'
NOMSG      DB 'NO FILE TRANSFER. RETURNING TO CPM.$'

```

```

ABRTMSG DB 'XMITING MICRO ABOFTED FILE TRANSFER.',13,10
DB 'PLEASE ERASE FILENAME FROM YOUR DIRECTORY.$'
FULLMSG DB 'DISK FULL. FILE TRANSFER INCOMPLETE.$'
ERMSG DB 'ENTER FILENAME AGAIN. END WITH <CR>$'
RXING1 DE 'CONNECTION MADE.$'
CONBUF DB 16H ;BUFFER FOR FILENAME
DB 00
DB 16H DUP(0)
BAUDMSG DB 'SELECT BAUD RATE',0DH,0AH
DB '1 = 300 BAUD',0DH,0AH
DB '2 = 600 BAUD',0DH,0AH
DB '3 = 1200 BAUD',0DH,0AH
DB '4 = 2400 BAUD',0DH,0AH
DB '5 = 4800 BAUD',0DH,0AH
DB '6 = 9600 BAUD$'
ERR1 DE 'BAUDRATE OUT OF RANGE$'
BAUD DW 00H
TABL DB 80H,01,0C0H,00,60H,00,3AH,00,18H,00,0CH,00
DATA ENDS
CODE SEGMENT PARA PUBLIC 'CODE'
START PROC FAR
    ASSUME CS:CODE
    PUSH DS
    MCV AX,0

```

```

PUSH AX
MOV AX,DATA
MOV ES,AX
ASSUME ES:DATA
MOV SI,80H
MOV DI,OFFSET DTA
MOV CX,80H
REP MOVSB
MOV SI,005CH
MOV DI,OFFSET FCB
MOV CX,12
REP MOVSB
MOV DS,AX
ASSUME DS:DATA
MOV DX,OFFSET BAUDMSG
MOV AH,09H
CALL RDOS
MOV AH,01H
CALL BDOS
SUB AL,31H
CMP AL,05H
JEE SETB1
JMP ERROR1

SETB1: MOV BX,OFFSET TABL
;TRANSFER BOTH PARAMETER AREAS TO OUR SGMMENT
;BAUDRATE HEADER
;PRINT SAME
;GET KEYBOARD INPUT
;CCNVERT TO TABLE OFFSET

```



```

ADD AL,AL
MOV AH,0
ADD BX,AX
MOV DX,[BX]
MOV BX,OFFSET BAUD
MOV [BX],DX
MOV DX,03FBH
MOV AL,83H
OUT DX,AL
MOV DX,03F8H
MOV BX,OFFSET BAUD
MOV AX,[BX]
OUT DX,AX
MOV DX,03FBH
MOV AL,03H
OUT DX,AL
MOV AH,1AH
MOV DX,OFFSET DTA
CALL BDOS

INIT: CALL CRLF
      MOV AH,09H
      MOV DX,OFFSET RIGHTS
      CALL BDOS
      CALL CRLF

      ;PRINT STRING TO SCREEN
      ;CCPYRIGHTS AND NAMES OF AUTHORS
      ;LINE CONTROL
      ;DIAB=1
      ;BAUDRATE DIVISOR
      ;CONTROL
      ;RESET DLAB
      ;OPEN THE DTA

```

```

CALL CRLF
MOV AH,09H
MOV DX,OFFSET WELCUM
CALL BDOS
CALL CRLF
CALL CRLF
MOV AH,09H
MOV DX,OFFSET INSTRC
CALL BDOS
CALL CRLF
CALL CRLF
HOLDING: MOV AH,06H
MOV DI,0FFH
CALL BDOS
AND AL,0DFH
CMP AL,53H
JNE G1
JMP MASTER
G1: CMP AL,52H
JNE G2
JMP SLAVE
G2: CMP AL,58H
JNE G3
JMP CPM
;PRINT STRING TO SCREEN
;WELCOME MSG
;PRINT STRING TO SCREEN
;SEND, RECEIVE, OR QUIT?
;CHECK FOR CONSOLE INPUT
;LOOKING FOR INPUT
;ENSURE LETTER IS A CAPITAL
;IS IT AN 'S'?
;IF SO, START FILE TRANSFER
;IS IT AN 'R'?
;IF SO, PREPARE TO RECEIVE FILE
;IS IT AN 'X'?
;IF SO, GO TO CPM

```

```

G3: JMP HOLDING          ;REPEAT UNTIL INPUT FOUND
MASTER: MOV AH,09H       ;PRINT STRING TO SCREEN
MOV DX,OFFSET ENTER     ;ENTER FILENAME
CALL BDOS
CALL CRLF

FILLUP: MOV BX,OFFSET FCB ;ADDRESS OF FCB
MOV [BX],BYTE PTR 00H
INC BX
MOV CL,0BH              ;11 SPACES
FILLUP1: MOV [BX],BYTE PTR 20H ;FILL MEMORY ADDRESS WITH SPACES
INC BX
DEC CL
JNZ FILLUP1
MOV CL,13H
FILLUP2: MOV [BX],BYTE PTR 00H ;FILL REST IF ADDRESS WITH 0'S
INC BX
DEC CL
JNZ FILLUP2
HOLD1: MOV AH,0AH
MOV DX,OFFSET CCNBUF    ;ADDRESS OF FIRST LETTER OF FILENAME
CALL BDOS
MOV BX,OFFSET CCNBUF    ;ADDRESS OF CONSOLE BUFFER
MOV DX,OFFSET FCB+1    ;FCB ADDRESS
INC BX

```

```

MOV CH,[BX]
MOV AL,[BX]
OR AL,AL
JNE G4
JMP ERROR
G4: INC BX
FLUP: MOV AL,[BX]
CMP AL,3AH
JNE G5
JMP DSKSEL
G5: CMP AL,2EH
JNE G6
JMP FIXIT
G6: CMP AL,40H
JNC G7
JMP DONTFIX
G7: AND AL,0DFH
DONTFIX: XCHG BX,DX
MOV [BX],AL
XCHG BX,DX
INC DX
INC BX
DEC CH
JNZ FLUP

;STORE COUNT IN BX REGISTER
;MCVE COUNT TO ACCUMULATOR
;IS THERE AN INPUT?

;TRY AGAIN

;IS CHARACTER A ':'?

;IF SO, GO TO DISK SELECT
;IS IT A '.'?

;IF SO, SKIP TO FILETYPE
;CHECK FOR LETTER

;SKIP NEXT STEP IF NOT LETTER
;ENSURE LETTER IS A CAPITAL

;STORE LETTER IN FCB

;REPEAT UNTIL END

```

```

LISN1: MOV AL,ATTN          ;LETTER 'R'
      CALL POUT1
      MOV CL,03H
LISN: CALL PIN
      CMP AL,RXACK
      JNE G8
      JMP XMIT
G8: DEC CL
      JNZ LISN
      JMP LISN1
XMIT: CALL CRLF
      CALL CRLF
      MOV AH,09H
      MOV DX,OFFSET RXING1
      CALL BDOS
      CALL CRLF
XMIT1: MOV AL,TXSYM
      CALL POUT1
      MOV CL,08AH
LIT1: CALL PIN
      CMP AL,TXACK
      JNE G9
      JMP TYFCB
G9: DEC CL
      ;IF SO, XMIT FILE CTRL BLK
      ;OTHERWISE KEEP LISTENING
      ;DC2 SYMBOL FOR SYNC AT START
      ;OF 128 BYTE BLOCK
      ;LISTEN 138 TIMES
      ;WAS 't' RECEIVED?

```

```

JNZ LITTLE
JMP XMIT1
TXFCB: CALL CRLF
CALL OPENIT
TXFCB1: MOV CH,00H
MOV BX,OFFSET FCB
FCBLUP: MOV AL,CH
INC BX
XOR AL,[BX]
MOV CH,AL
MOV AL,[BX]
CALL POUT1
CMP AL,0H
JNE G10
JMP FCBCX
G10: JMP FCBLUP
FCBCX: MOV CL,20H
FCBCX1: CALL PIN
CMP AL,RXACK
JNZ FCBCX1
MOV AL,CH
CALL POUT
PUSH CX
MOV AL,0H
;UNTIL CTR IS ZERO,
;THEN SEND DC2 SYNC AGAIN
;SEND FILENAME TO RXING MICRO
;SEE IF FILE EXISTS. IF SO, OPEN IT
;INITIALIZE CHECKSUM REGISTER
;SET PTR TO 1ST LETTER IN FILENAME
;PERFORM CHECKSUM OPERATION
;MOVE PTR TO NEXT BYTE
;BY XORING CURRENT BYTE
;WITH B REGISTER
;PUT CURRENT BYTE IN ACCUM
;SEND CURRENT BYTE
;CHECK FOR END OF FILENAME
;IF END, GO TO CHECKSUM LOOP
;IF NOT, REPEAT FCB LOOP
;LCOP 32 TIMES
;FCR SYNC WITH SLAVE
;IS IT AN 'I'?
;IF NOT, LISTEN AGAIN. IF SO,
;PUT CHECKSUM IN ACCUM
;SEND CHECKSUM
;SAVE CHECKSUM
;CLEAR ACCUM

```



```

MOV      CH,80H
FCB TMCT: CALL PIN
CMP      AL,BAD
JNE G11
JMP      RSNDFCB
G11:    CMP      AL,GOOD
JNE G12
JMP      WAITFIL
G12:    DEC      CH
JNZ      FCB TMOT
POP      CX
DEC      CL
JNZ      FCBCK1
JMP      TXFCB
WAITFIL: POP      CX
WAIT2:  MOV      CX,07FFH
WAIT1:  CALL     STATIN1
JZ       WAIT1
DEC      CX
JNE G13
JMP      GOCPM
G13:    CALL     PIN
CMP      AL,QUIT
JNE G14

;LISTEN 100 TIMES
;READ MAIL
;DID IT CHECK BAD?

;IF SO, SEND FCB AGAIN
;DID IT CHECK GOOD?

;IF SO, GO TO NEXT ROUTINE
;IF NOT, DECREMENT CTR, AND
;IF NOT 0, LISTEN AGAIN
;CLEAR STACK
;IF SO, DECREMENT CL
;AND REPEAT UNTIL CL=0
;IF 0, ASSUME PROBLEM AND SEND AGAIN
;CLEAR STACK
;CCUNT LOOP APPX 2K
;ANY 'MAIL'?
;IF NOT, CHECK AGAIN
;IF SO, DECREMENT CTR

;AND, IF 0, QUIT
;OTHERWISE READ 'MAIL'
;DCES RXING MICRO AIFREADY HAVE FILE?

```

```

JMP GOCPM1
G14: CMP AL,GOON
JNZ WAIT2
CALL CRLF
TXDATA: MOV AH,09H
MOV DX,OFFSET TXING1
CALL BDCS
CALL CRLF
RDSEQ: CALL READSEQ
SEND: CALL CHECK
MOV AL,TXSYH
CALL POUT1
MOV CL,0FH
LITLET2: CALL PIN
CMP AL,TXACK
JNE G15
JMP SLUP2
G15: CMP AL,DSKFUL
JNE G16
JMP FULDISK
G16: DEC CL
JNZ LITLET2
JMP SEND
;IF SO, GO TO CPM
;IS IT THE GO ON SIGNAL 'G'
;IF NOT, LISTEN AGAIN. ALLOW RXING
;MICRO TO CATCH UP
;PRINT STRING TO SCREEN
;SAYS FILE BEING SENT
;READ FIRST 128 BYTE BLOCK
;AND SEND TO RXING MICRO
;DC2 SYMBOL FOR SYNC AT START OF DATA
;LISTEN 15 TIMES
;IS IT A 't'?
;IF SO, READY TO SEND DATA
;IS RXING MICRO'S DISK FULL?
;IF SO, QUIT
;IF NOT, DECREMENT CTR
;LISTEN AGAIN, UNLESS CTR IS 0,
;THEN TRY TO SYNC AGAIN

```

```

SLUP2: MOV AL,RLDTA
      CALL POUT1
      MOV CX,07FFH
SLUP3: CALL PIN
      CMP AL,RLDTA
      JZ SLUP1
      DEC CX
      JNZ SLUP3
      JMP SLUP2

SLUP1: MCV BX,OFFSET DTA
      MOV CH,00H
      MCV CL,80H
      SLOOP: MOV AL,CH
      XCR AL,[BX]
      MOV CH,AL
      MCV AL,[BX]
      CALL POUT
      INC BX
      DEC CL
      JNZ SLOOP

CRC: MOV AL,CH
      CALL POUT
CRCTMCT: MOV CH,01AH
CRCT1: CALL STATIN1
;OCBH MEANS TIME FOR DATA
;WAIT LOOP APPX 2K
;IS IT ECHO?
;IF SO, SEND DATA
;DECREMENT COUNTER
;REPEAT UNTIL ZERO
;PCINTER TO 1ST INFO BYTE
;INITIALIZE CHECKSUM LOCATION
;PERFORM CHECKSUM
;XCR DATA WITH CH REGISTER
;PUT BYTE IN ACCUMULATOR
;DATA IS TRANSFERRED
;MCVE PTR TO NEXT BYTE
;PUT CHECKSUM IN ACCUMULATOR
;AND SEND TO RXING MICRO
;LISTEN 26 TIMES
;CHECK INPUT BUFFER

```

```

JZ   CRCT1
CALL PIN
CMP  AL,BAD
JNE G17
JMP RESEND
G17: CMP AL,GOOD
JNE G18
JMP EDSQRPT
G18: DEC CH
JNZ CRCT1
JMP SEND
DSKSEL: MOV BX,OFFSET CONBUF+2 ;ADDRESS OF DISK SEL ENTRY
MCV AL,[BX]
AND AL,0DFH
CMP AL,'A'
JZ ALISK
CMP AL,'B'
JZ BDISK
CMP AL,'C'
JZ CDISK
JMP DSKSEL1
ADISK: MOV DI,OFFSET FCB
MOV [DI],BYTE PTR 01H
;IF NOTHING, TRY AGAIN
;READ MAIL
;IS CHECK BAD?
;IF SO, SEND BLOCK AGAIN
;IS CHECK GOOD?
;IF SO, READ NEXT BLOCK
;DECREMENT COUNTER
;IF NOT TIMED OUT,LISTEN AGAIN
;IF TIMED OUT, ASSUME PROBLEM.
;SEND BLOCK AGAIN
;PUT DISK SEL IN ACCUM
;ENSURE LETTER IS CAPITAL
;IS LETTER AN 'A'?
;IF SO, SET FOR A DRIVE.
;IS LETTER A 'B'?
;IF SO, SET FOR B DRIVE.
;IS LETTER A 'C'?
;IF SO, SET FOR C DRIVE.
;IF NEITHER, RETURN TO FILENAME LOOP.
;SET PTR TO DRIVE BYTL.
;SET FCB FOR A DRIVE.

```

```

JMP DSKSEL1
BDISK: MOV DI,OFFSET FCB
MOV [DI],BYTE PTR 02H
JMP DSKSEL1
CDISK: MOV DI,OFFSET FCB
MOV [DI],BYTE PTR 03H
DSKSEL1: INC BX
INC PX
MOV DX,OFFSET FCB+1
JMP FLUP
RESEND: MOV AH,CONCUT
MOV DL,BAD
CALL BDOS
JMP SEND
RSNDFCB: POP CX
MOV AL,CH
ADD AL,3
CALL POUT1
RSNDF: CALL PIN
CMP AL,RXACK
JNZ RSNDF
RSNDFC1: CALL PIN
CMP AL,TXACK
JNZ RSNDFC1
;RETURN TO FILENAME LOOP.
;SET PTR TO DRIVE BYTE.
;SET FCB FOR B DRIVE.
;RETURN TO FILENAME LOOP.
;SET PTR TO DRIVE BYTE.
;SET FCB FOR C DRIVE.
;MOVE BUFFER POINTER TO FILENAME.

;FCB FILENAME ADDRESS.
;RETURN TO FILENAME LOOP.
;PINT TO SCREEN
;A 'b' IF CHECKSUM WAS BAD

;AND SEND BLOCK AGAIN
;RFCALL CHECKSUM
;PUT CHECKSUM IN ACCUM
;ADD 3 TO OFFSET
;SEND BYTE

;IS IT AN 'r'
;IF NOT LISTEN AGAIN
;READ MAILBOX
;SYNC WITH RXING MICRO
;REPEAT UNTIL TXACK RECEIVED

```

```

JMP TXFCB1
;IF SO, RESEND FCB
FIXIT: MOV DX,OFFSET FCB+9 ;MCVE POINTER TO FILETYPE AREA
INC BX ;MCVE PTR TO FIRST LETTER OF FILTYPE
JMP FLUP
ERROR: MOV AH,09H ;PFINT STRING TO SCREEN
MOV DX,OFFSET ERMMSG ;ERROR MESSAGE
CALL BDOS
CALL CRLF
JMP HOLD1
;LCOK FOR INPUT AGAIN
POUT1 PROC NEAR
LAHF
PUSH AX
;SEND THE DATA
;FIRST, SAVE THE CURRENT BYTE
MOV AH,06H ;CHECK FOR CONSOLE INPUT
MOV DL,0FFH ;LCOKING FOR INPUT
CALL BDOS
CMP AL,CTRLC ;IS THERE A CONTROL C?
JNE G22
JMP STOPS
G22: CALL CHECK
PCP AX
;PERFORM CHECK
SAHF ;AND RECALL BYTE
MOV DX,DATA1
OUT DX,AL

```



```

RET
POUT1 ENDP
OPENIT PROC NEAR
    MOV     AH,0FH
    MOV     DX,OFFSET FCB
    CALL    BDOS
    CMP     AL,0FFH
    JNE     G23
    JMP     FNFOUND
G23: RET
OPENIT ENDP
CLOSEIT: MOV     AH,10H
    MOV     DX,OFFSET FCB
    CALL    BDOS
    CLOSEIT1: MOV     AL,DONE
    CALL    POUT1
    MOV     AL,0H
    CALL    PIN
    CMP     AL,DONE
    JNZ     CLOSIT1
    JMP     GOCPM
READSEQ PROC NEAR
    PUSH    CX
    PUSH    DX
;OPEN FILE CODE
;FILE CTRL BLOCK ADDRESS IN LX REG PR
;FF = FILE NOT FOUND
;IF FILE NOT FOUND
;OTHERWISE, RET TO TX DATA
;CLOSE FILE CODE
;FILE CTRL BLOCK ADDRESS IN DE REGPR
;END OF SESSION MSG 'Z'
;SEND TO RXING MICRO
;CLEAR ACCUM
;CHECK REPLY
;DCES RXING MICRO AGREE?
;IF NOT, REPEAT
;IF SO, GO TO CPM

```

```

MOV     AH,14H
MCV     DX,OFFSET FCB
CALL    BDOS
POP     DX
POP     CX
CMP     AL,0
JZ      G24
JMP     EOFILE
G24:    RET

READSEQ ENDP
RDSQRPT: MOV  AH,CONCUT
MOV     DL,02AH
CALL    BDOS
JMP     RDSEQ
FNFOUND: MOV  AL,QUIT
CALL    POUT1
MOV     AH,09H
MCV     DX,OFFSET FNFDMSG
CALL    BDOS
CALL    CRLF
JMP     GOCPM
EOFILE: POP  AX
EOFILE2: MOV  AL,TSYMS
CALL    POUT1

;READ SEQUENTIAL CODE
;FILE CTRL BLOCK ADDRESS IN DX REGPR

;0 MEANS SUCCESSFUL READ
;IF NOT 0, ASSUME FINISHED WITH FILE

;PRINT TO SCREEN
;*' SO USER KNOWS BLK WAS SENT

;TC READ NEXT 128 BYTE BLK
;TELL RXING MICRO NO FILE FOUND

;PRINT STRING TO SCREEN
;FILE NOT FOUND MSG

;AND GO TO CPM
;CCRRRECT STACK POINTER
;DC2 SYMBOL FOR SYNC WITH FXING MICRO

```

```

MOV     CL,0FH
LITLET3: CALL STATIN1
        JZ     LITLET3
        CALL PIN
        CMP    AL, TXACK
        JNE    G25
        JMP    EOFIL1
G25:    DEC    CL
        JNZ    LITLET3
        JMP    EOFILE2
EOFIL1: MOV    AL, QUIT
        CALL POUT1
        CALL PIN
        CMP    AL, QUIT
        JNZ    EOFIL1
        CALL CRLF
        MOV    AH, 09H
        MOV    DX, OFFSET EOFMSG
        CALL BDOS
        CALL CRLF
        JMP    CLOSIT
STOPS:  MOV    AL, CTRLC
        MOV    EX, DATA1
        OUT    EX, AL

;LISTEN 15 TIMES
;CHECK FOR MAIL
;IF NONE, CHECK AGAIN
;READ MAIL
;IS IT A 't'?
;IF SO, CONTINUE
;IF NOT, DECREMENT COUNTER
;AND LISTEN AGAIN, UNLESS COUNTER IS
;0. THEN TRY AGAIN
;DC4 SYMBOL. TELLS RXING MICRO THAT
;THE FILE IS DONE
;LISTEN FOR REPLY
;DCES RXING MICRO ACKNOWLEDGE?
;IF NOT, TRY AGAIN
;PRINT STRING TO SCREEN
;IF SO, TELL USER FILE IS DONE
;AND CLOSE THE FILE
;SEND CTRLC TO RXING MICRO

```

```

MOV AL,0H
CALL PIN
CMP AL,CTRLC
JNZ STOPS
POP AX
JME GOCPM
FULDISK: MOV AL,DONE
CALL POUT1
MOV AH,09H
MOV DX,OFFSET FULMSG
CALL BDOS
CALL CRLF
JMP GOCPM
GOCPM: MOV AL,0H
MOV DX,DATA1
OUT DX,AL
CALL CRLF
CALL PIN
JMP CPM
GOCPM1: MOV AH,09H
MOV DX,OFFSET HASFILE
CALL RDOS
JME GOCPM
SLAVE: MOV AH,09H
;CLEAR ACCUM
;FROM RXING MICRO
;ACK FROM RXING MICRO
;REPEAT UNTIL ACK
;LETTER 'Z' TO ACKNOWLEDGE
;SEND BYTE
;PRINT STRING TO SCREEN
;SAYS RXER'S DISK FULL
;RESET THE ACCUMULATOR AND
;CLEAR OUTPUT BUFFER
;AND GO TO CPM
;PRINT STRING TO SCREEN
;RXING MICRO HAS FILE ALREADY
;PRINT STRING TO SCREEN

```

```

MOV DX,OFFSET WCHDSK ;SELECT DISK DRIVE
CALL BDOS
CALL CRLF
CALL CRLF

DRVSEL: MOV AH,06H ;CHECK FOR CONSOLE INPUT
        MOV DL,0FFH ;ICOKING FOR INPUT
CALL BDOS
CMP AL,0DH ;IS IT A <CR>?
JNE G26
JMP CONT ;IF SO, ENTER RECEIVE MODE

G26: AND AL,0DFH ;ENSURE LETTER IS A CAPITAL
CMP AL,'A' ;IS IT AN 'A'?
JZ G27
JMP DISKB ;SKIP TO B IF NOT 'A'

G27: MOV BX,OFFSET FCB ;ADDRESS OF DISK DRIVE BYTE
MOV [BX],BYTE PTR 01H ;SET BYTE TO A DISK DRIVE
JMP CONT ;THEN CONTINUE

DISKB: CMP AL,'B' ;IS IT A 'B'?
JNZ DISKB ;SKIP TO C IF NOT 'B'

MOV BX,OFFSET FCB ;ADDRESS OF DISK DRIVE BYTE
MOV [BX],BYTE PTR 02H ;SET BYTE TO B DISK DRIVE
JMP CONT ;THEN CONTINUE

DISKC: CMP AL,'C' ;IS IT A 'C'?
JNZ DRVSEL ;IF NOT, LISTEN AGAIN

```

```

MOV    BX,OFFSET FCB          ;ADDRESS OF DISK DRIVE BYTE
MOV    [BX],BYTE PTR 03H      ;SET BYTE TO C DISK DRIVE
CONT:  MOV    AH,09H           ;PRINT STRING TO SCREEN
MOV    DX,OFFSET RXMODE       ;IN RECEIVE MODE
CALL  BDOS
CALL  CRLF

SLAVE1: MOV    AL,00H         ;RESET ACCUMULATOR
CALL  PIN                    ;LISTENING FOR AN 'R'
CMP    AL,ATTN               ;'F'
JNZ    SLAVE1                ;IF 'R' RX'D, CONTINUE. IF NOT
                                   ;LISTEN AGAIN

CALL  CRLF
MOV    AH,09H                ;PRINT STRING TO SCREEN
MOV    DX,OFFSET RXING1      ;CCNECTION MADE
CALL  BDOS
CALL  CRLF
MOV    AL,RXACK
CALL  POUT
LISTEN: CALL PIN1
CMP    AL,TSYM
JNZ    LISTEN
                                   ;'I'
                                   ;SEND AN 'I' TO XMITING MICRO
                                   ;LISTENING FOR A 'DC2'
                                   ;'IC2'
                                   ;IF 'DC2' KX'D, CONTINUE. IF NOT,
                                   ;LISTEN AGAIN

CALL  CRLF

```

```

RXFCB: MOV    BX,OFFSET FCB+1 ;ADDRESS OF FCB MEM LOC INTO BX REG PR

```



```

MCV CL,1EH ;CCUNTER FOR FCB'S 31 SPACES
RSTFCF: MOV [BX],BYTE PTR 00H ;FILL FCB WITH )'S
INC BX ;MOVE PTR TO NEXT MEMORY ADDRESS IN FCB
DEC CL ;DECREMENT COUNTER
MOV AL,CL
CMP AL,0
JNZ RSTFCB
;IF COUNTER = 0, CONT. IF NOT,
;PUT ANOTHER 0 IN FCB
;INITIALIZE CHECKSUM
MOV CH,00H
MCV BX,OFFSET FCB+1
MOV AL,TXACK
CALL POUT
CALL PIN1
RST1: CALL STATIN1
JZ RST1
RST2: CALL PIN1
CMP AL,QUIT
JNE G28
JMP NOFILE
G28: CMP AL,0H
JNE G29
JMP FCBCRC
G29: CMP AL,TXSYM
JZ RST2
;FILE NAME DATA
;IS DATA A 'QUIT'?
;FILE DID NOT EXIST
;CHECK IF FILENAME COMPLETELY SENT
;IF FILENAME RX'D, GO TO CHECKSUM
;CHECK IF DATA IS VALID
;IF DATA IS NOT FILENAME,

```

```

MOV [BX],AL
CALL OUTPUT
MOV AL,CH
XOR AL,[BX]
MOV CH,AL
INC BX
JMP RST1

FCBCRC: MOV AL,FXACK
CALL POUT
FCBCRC1: CALL STATIN1
JZ FCBCRC1
CALL CR1F
CALL PIN
CMP AL,CH
JNE G30
JMP STRTFIL
G30: ADD AL,3
MOV CL,AL
MOV AL,BAD
CALL POUT
CLEAR: CALL PIN
CMP AL,CL
JNZ CLEAR
MOV AL,FXACK

;PUT FILENAME IN FCB
;PRINT FILENAME TO SCREEN
;CALCULATE CHECKSUM
;MOVE PTR TO NEXT FCB ADDRESS
;'I'
;SYNC DATA WITH XMITTING MICRO
;CHECKING FOR INPUT
;CHECKSUM DATA
;COMPARE CHECKSUM
;CHECKSUM MATCHED
;ADD 3 TO THE CHECKSUM
;STORE IN REGISTER
;CHECKSUM DID NOT MATCH
;TELL XMITTING MICRO
;XMITTING MICRO STOPPED SENDING CHECKSUM?
;IF NOT,LISTEN AGAIN
;SYNC WITH XMITTING MICRO

```

```

CALL POUT
JMP RXFCB
STRTFIL: MOV AL,GOOD
CALL POUT
CALL OPNFILE
CALL MAKEFIL
RXD1: MOV CH,00H
MOV BX,OFFSET DTA
MOV CL,81H
RXDS: CALL STATIN1
JZ RXDS
RXD2: CALL PIN1
CMP AL,TXSYM
JNZ RXD2
MOV AL,TXACK
CALL POUT
RXDS1: CALL STATIN1
JZ RXDS1
RXDET: CALL PIN1
CMP AL,R1DTA
JZ RXYET1
CMP AL,QUIT
JNZ RXYET4
;TRY AGAIN
;READY TO CHECK IF FILE ALREADY PRESENT
;CHECK IF FILE EXISTS
;CREATE NEW FILE
;INITIALIZE CHECKSUM
;ICAD ADDRESS OF DMA MEM LOC TO
;BX REGISTER PAIR
;INITIALIZE COUNTER WITH SIZE OF DMA
;SYNC WITH XMITING MICRO
;COMPARE WITH 'DC2'
;'T'
;IN SYNC WITH XMITING MICRO
;CHECKING FOR INPUT
;IS IT OCBH?
;IF SO, GO TO RECEIVE DATA
;IS IT 'DC4' FOR QUIT?
;IF NOT, LISTEN AGAIN; OTHERWISE,

```

```

JMP CLSFILE
RXYET4: JMP RXDS1
RXYET1: MOV AL,RLDTA
        CALL POUT
        MOV AL,00H
RXYET2: CALL STATIN1
        JZ RXYET2
RXYET3: CALL PIN
        CMP AL,RLDTA
        JZ RXYET3
RXD3: DEC CL
        JNE G32
        JMP RXCRC
G32: MOV [BX],AL
        MOV AL,CH
        XCR AL,[BX]
        MOV CH,AL
        INC BX
RXD4: CALL STATIN1
        JZ RXD4
        CALL PIN
        JMP RXD3
RXCRC: MOV AL,AL
        CMP AL,CH
;IF SO, CLOSE FILE
;LISTEN AGAIN
;ACK REAL DATA COMING
;CLEAR ACCUM
;CHECKING FOR INPUT
;READ DATA
;IS IT STILL RLDTA?
;DECREMENT COUNTER
;CHECKSUM RX'D
;PUT THE DATA IN MEMORY
;CALCULATE CHECKSUM
;MCVE PTR TO NEXT DMA ADDRESS
;CHECK FOR INPUT
;LCOP UNTIL INPUT
;ENSURE CH IS COMPARED TO A
;CCMPARE WITH CHECKSUM

```

```

JNE G33
JMP WRITFIL
G33: MOV AL,BAD
CALL POUT
JMP RXD2
POUT PRCC NEAR
LAHF
PUSH AX
CALL CHECK
POP AX
SAHF
MOV DX,DATA1
OUT DX,AL
RET

POUT ENDP
PIN1 PROC NEAR
MOV IX,DATA1
IN AL,DX
CMP AL,CTRLC
JNE G34
JMP ABORT
G34: RET
PIN1 ENDP
WRITFIL: MOV AL,GOOD
;128 BYTE BLOCK SENT
;CHECKSUM DID NOT MATCH
;NCTIFY XMITING MICRO
;SEND 128 BYTE BLOCK AGAIN

;SAVE THE DATA

;RETURN THE DATA

;SEND DATA

;DID XMITING MICRO ABORT?
;IF SO, ABORT

;XMIT THAT THE CHECKSUM IS CORRECT

```

```

CALL POUT
CALL WRITSEQ
MOV AH,CONOUT
MOV DL,02AH
CALL BDOS
JMP RXD1
OPNFILE PROC NEAR
MOV AH,0FH
MOV DX,OFFSET FCB
CALL BDOS
CMP AL,0FFH
JZ G35
JMP FILFND
G35: RET
OPNFILE ENDP
CLSFILE: MOV AL,QUIT
CALL POUT
MOV AH,10H
MOV DX,OFFSET FCB
CALL BDOS
CALL CRLF
MOV DX,OFFSET ECFMSG
MOV AH,09H
CALL BDOS
;START WRITING FILE TO DISK
;PRINT TO SCREEN
;*' TO PRINT TO SCREEN
;OEEN FILE CODE
;FCB ADDRESS IN DX RGSTR PAIR
;FF = FILE NOT FOUND
;FILE EXISTS
;'IC4'
;AGREE END OF FILE
;CIOSE FILE CODE
;FCB ADDRESS IN DX RGSTR PAIR
;FILE TRANSMISSION COMPLETED
;PRINT STRING TO SCREEN

```


CALL CRLF	
MOV AL,0H	;CLEAR THE ACCUMULATOR
CLSFIL1: CALL PIN1	;LOOKING FOR END OF SESSION MSG
CMP AL,DONE	; 'Z' = END OF SESSION
JNZ CLSFIL1	
MOV AL,DONE	;END OF SESSION MESSAGE
CALL POUT	;CCNFIRM RECEPTION OF E-O-SESSION MSG
JME CPM	
NAKEFIL PROC NEAR	
MOV AH,16H	;MAKE NEW FILE CODE
MOV DX,OFFSET FCB	;FCB ADDRESS IN DX RGSTR PAIR
CALL BDOS	
MOV AL,GOON	;CCNTINUE MESSAGE
CALL POUT	
RET	;RETURN TO RX FIRST 128 BYTE BLOCK
MAKEFIL ENDP	
WAITSEQ PROC NEAR	
PUSH CX	
PUSH DX	
MOV AH,15H	;WRITE THE FILE TO THE DISK
MOV DX,OFFSET FCB	;FCB IN DX RGSTR PAIR
CALL BDOS	
POP DX	
POP CX	

```

OR AL,AL
JZ G36
JMP FULLDSK
G36: RET
WRITEQ ENDP
FILEND: MOV AL,QUIT
CALL POUT
MOV AH,09H
MOV DX,OFFSET FNDMSG
CALL BDOS
CALL CRLF
JMF CPM
NOFILE: MOV AH,09H
MOV DX,OFFSET NCMSG
CALL BDOS
CALL CRLF
JMF CPM
ABORT: CALL CRLF
MOV AL,CTPLC
CALL POUT
MOV AH,09H
MOV DX,OFFSET AERTMSG
CALL BDOS
CALL CRLF
;CHECK IF DISK IS FULL
;IF SO, JUMP TO FULLDSK
;TELL XMITING MICRO, FILE FOUND
;PRINT STRING TO SCREEN
;FILE ALREADY EXISTS. GO TO CPM
;PRINT STRING TO SCREEN
;NC FILE TRANSFER
;SEND XMITING MICRO ABORT ACK
;PRINT STRING TO SCREEN
;XMITING MICRO ABORTED

```

```

JMF GOCPM
FULLDSK: MOV AL,DSKFUL
CALL POUT
CALL PIN
CMP AL,DONE
JNZ FULLDSK
MOV AH,09H
MOV DX,OFFSET FULLMSG
CALL BDOS
CALL CRLF
JMF CPM

CHECK PROC NEAR
CALL STATIN2
JZ CHECK
RET

CHECK ENDP
STATIN1 PROC NEAR
MOV DX,STATUS
IN AI,DX
AND AL,EXEDY
RET
STATIN1 ENDP
STATIN2 PROC NEAR
MOV DX,STATUS
;GC TO CPM
;I'
;TELL XMITING MICRO DISK FULL
;AWAITING CONFIRMATION
;PRINT TO SCREEN
;FILE TRANSFER INCOMPLETE, DISK FULL
;CHECK STATUS BYTE
;CONTINUE UNTIL TXRDY IS SET

```

```

IN AL,DX
AND AL,TXRDY
RET
STATIN2 ENDP
PIN PROC NEAR
    MOV DX,DATA1
IN AL,DX
RET
PIN ENDP
CRLF PROC NEAR
    PUSH AX
    MOV AL,0DH
    CALL OUTPUT
    MOV AL,0AH
    CALL OUTPUT
    POP AX
    RET
CRLF ENDP
OUTPUT PROC NEAR
    PUSH ES
    PUSH EX
    PUSH DX
    PUSH CX
    LAHF
    ;SAVE THE ES,
    ;EX
    ;DX,
    ;AND CX REGISTERS

```

```

PUSH AX
MOV AH,CONOUT
MCV DL,AL
CALL BDOS
PCP AX
SAHF
POP CX
POP DX
PCP BX
POP ES
RET
1 OUTPUT ENDP
4 9 BDOS PROC NEAR
PUSH ES
PUSH BX
PUSH DX
PUSH CX
INT 21H
POP CX
POP DX
PCP BX
PCP ES
RET
BDOS ENDP
;PRINT TO SCREEN
;PUT THE ACCUMULATOR IN 'DI' RGSTR
;RETURN THE CX,
;DX,
;BX,
;AND ES REGISTERS
;SAVE THE ES
;BX,
;DX,
;AND CX REGISTERS
;EXECUTE
;RETURN THE CX,
;DX,
;BX,
;AND ES REGISTERS

```

```

CPM: RET
ERROR1: MOV DX,OFFSET ERR1
        MOV AH,09H
        CALL BDOS
        RET
START ENDP
CODE ENDS
END START
;FAR RETURN
;BAUDRATE ERROR
;PRINT STRING TO SCREEN
;FAR RETURN

```

NOTE: This version of MICROLAN was obtained by converting the CP/M-86 version to meet the requirements of the MS-DOS operating system. We did not attempt to streamline the procedure or efficiency of operation for this IBM compatible version.

APPENDIX F
MICROLAN'S ON SCREEN MESSAGES

BAUDMSG	SELECT BAUD RATE
	1 = 300 BAUD
	2 = 600 BAUD
	3 = 1200 BAUD
	4 = 2400 BAUD
	5 = 4800 BAUD
	6 = 9600 BAUD

ERR 1	BAUDRATE OUT OF RANGE
-------	-----------------------

RIGHTS	MICROLAN VERSION 2.1
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WELCUM	WELCOME!
	YOU ARE NOW ENTERING THE TRANSFER ZONE!

INSTRC	ENTER AN S FOR TRANSMIT MODE, AN R FOR RECEIVE MODE, OR AN X TO EXIT.
--------	--

ENTER	ENTER NAME OF FILE TO BE SENT. IF THE FILE IS ON A DISK IN ANOTHER DRIVE, ENTER IN THE FORMAT: B:FILENAME.FILETYPE
-------	--

ERMSG	ENTER FILENAME AGAIN. END WITH <CR>
-------	-------------------------------------

WCHDSK	WRITE FILE TO WHICH DISK DRIVE? ENTER AN A FOR A DRIVE, A B FOR B DRIVE,... OR PRESS RETURN FOR DEFAULT DRIVE.
EXMODE	IN RECEIVE MODE.
NOMSG	NO FILE TRANSFER. RETURNING TO CPM.
FNFDMSG	FILE DOES NOT EXIST, RETURNING TO CPM.
FNDMSG	FILE ALREADY EXISTS. RETURNING TO CPM.
HASFILE	RXING MICRO HAS FILE ALREADY, GOING TO CPM.
RXING1	CONNECTION MADE.
TXING1	TRANSMITTING FILE.
FULLMSG	DISK FULL. FILE TRANSFER INCOMPLETE.
FULMSG	RXING MICRO DISK FULL. RETURNING TO CPM.
ABRTMSG	XMITING MICRO ABORTED FILE TRANSFER. PLEASE ERASE FILENAME FROM YOUR DIRECTORY.
EOFMSG	FILE TRANSMISSION COMPLETED.

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